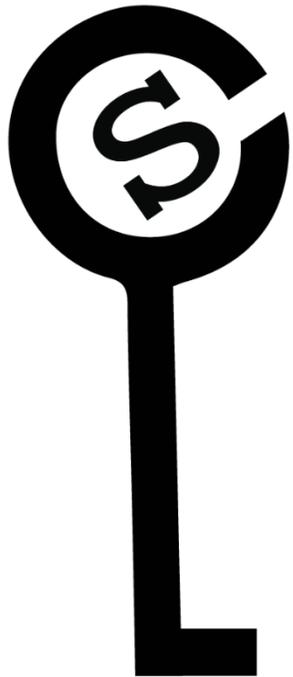


# The Interpose PUF (iPUF): Secure PUF Design against State-of-the-art Machine Learning based Modeling Attacks

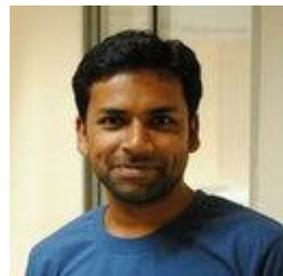
---



**Phuong Ha Nguyen,**  
**Durga P. Sahoo, Kaleel Mahmood, Chenglu Jin,**  
**Ulrich Rührmair and Marten van Dijk**



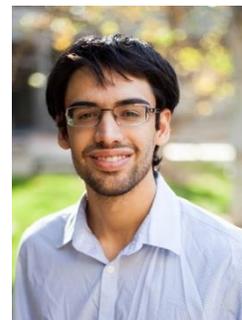
Ha



Durga



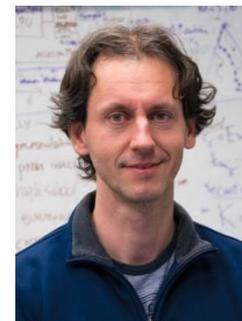
Chenglu



Kaleel



Uli



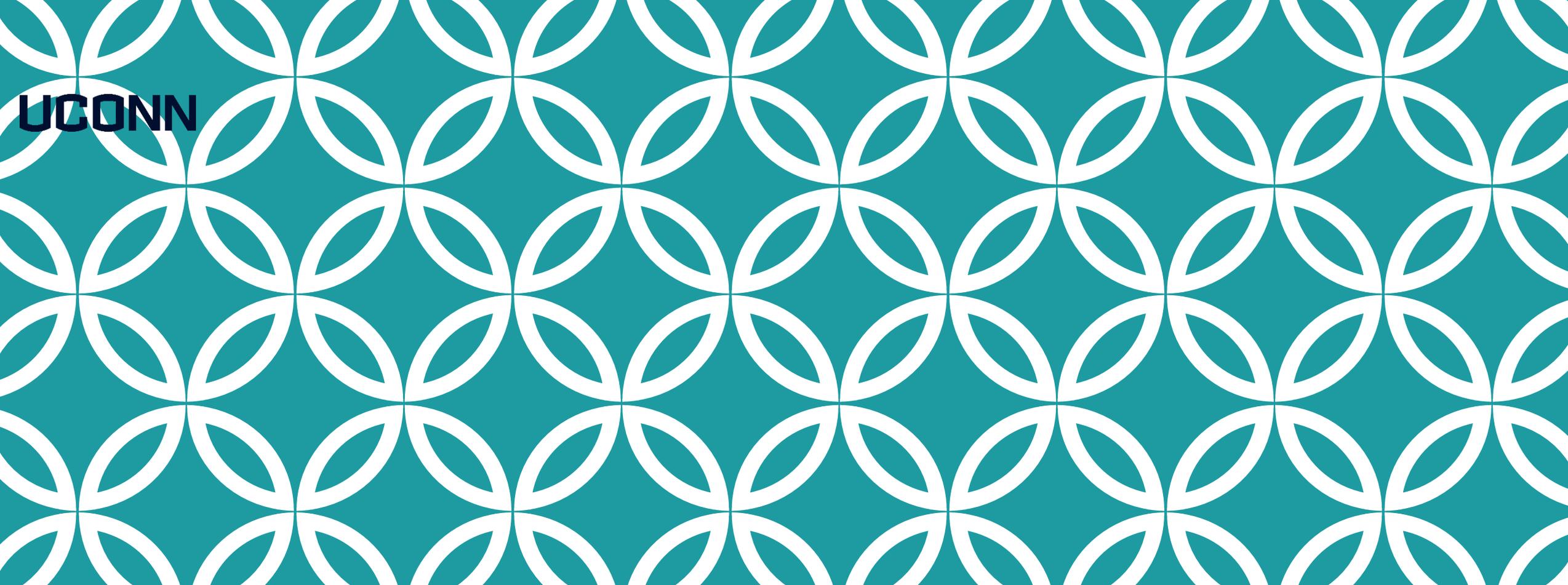
Marten

**CHES 2019**

**UCONN**

1. Concept - Overview - Motivation
2. Strong PUFs: APUF, XOR APUF and Interpose PUF (iPUF)
3. Short-term Reliability
4. Reliability based modeling attacks on XOR PUF: understanding
5. Interpose PUF – a lightweight PUF which is secure against state-of-the art modeling attacks
6. Conclusion





**UConn**

# 1. Concept - Overview - Motivation



# Concept - Overview – Motivation [1]

---



# Concept - Overview – Motivation [2]

---



**Nature:** process variation – physically unclonability - unique

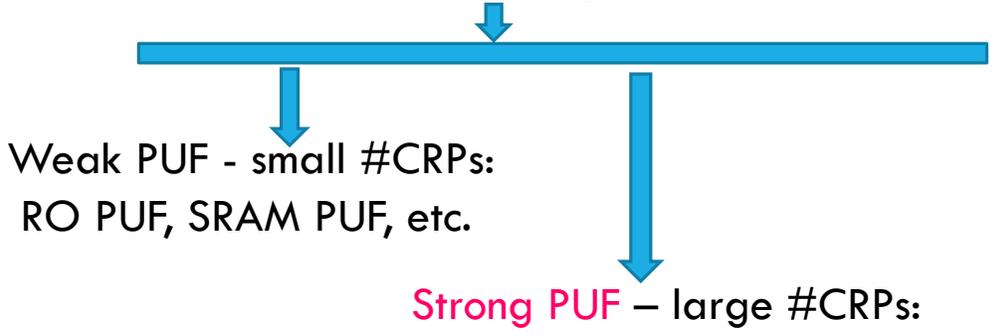
**Application:** device Identification, authentication  
and crypto key generation



# Concept - Overview – Motivation [3]



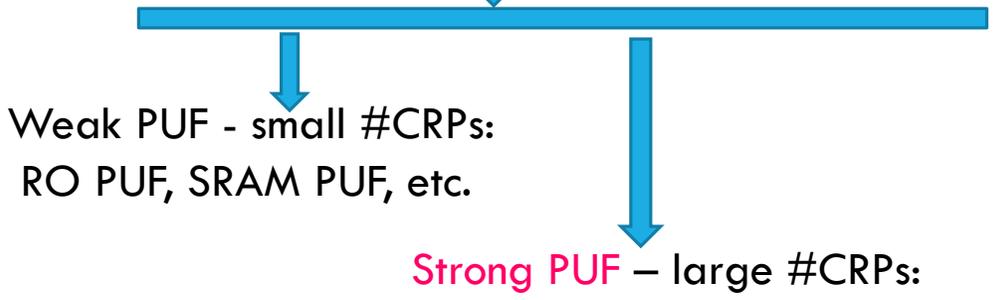
PUF's Category:



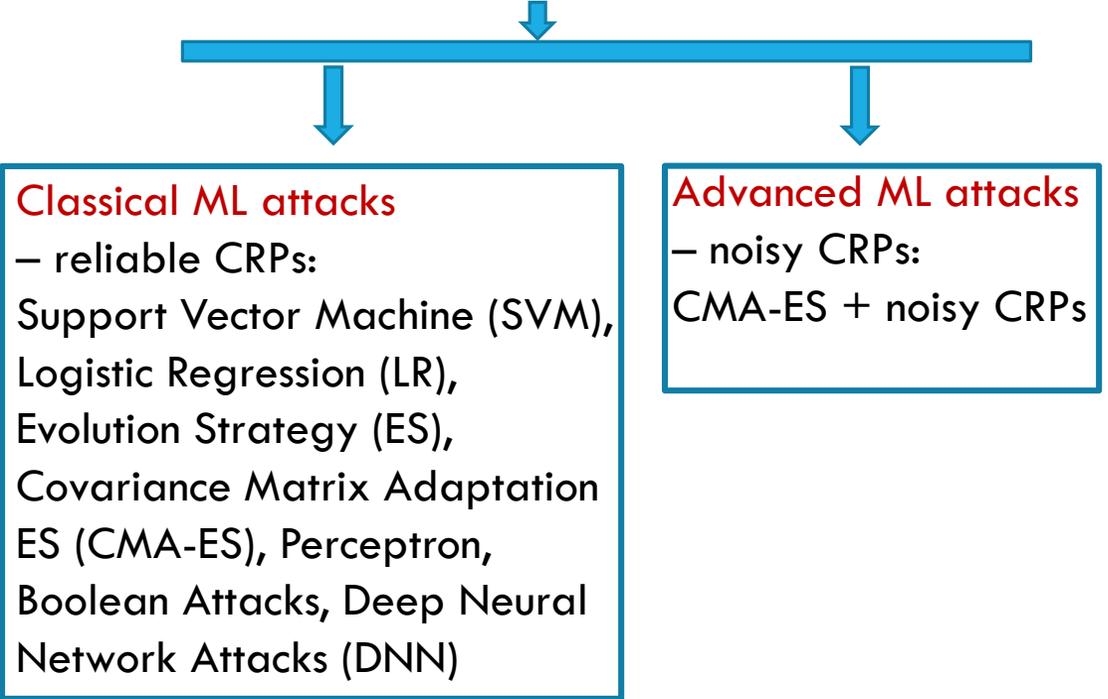
# Concept - Overview – Motivation [4]



PUF's Category:



PUF's Modeling Attacks with CRPs only:



# Concept - Overview – Motivation [5]



PUF's Category:

PUF's Modeling Attacks with CRPs only:

Weak PUF - small #CRPs:  
RO PUF, SRAM PUF, etc.

Strong PUF – large #CRPs:

**Broken but lightweight:**  
Arbiter PUF/APUF, XOR APUF, Feed Forward PUF, Lightweight Secure PUF, Bistable Ring PUF.

**No Security Proof:**  
Power Grid PUF, Clock PUF, Crossbar PUF

**Security Proof:**  
LPN PUFs - Large HW footprint

**Classical ML attacks**  
– reliable CRPs:  
Support Vector Machine (SVM), Logistic Regression (LR), Evolution Strategy (ES), Covariance Matrix Adaptation ES (CMA-ES), Perceptron, Boolean Attacks, Deep Neural Network Attacks (DNN)

**Advanced ML attacks**  
– noisy CRPs:  
CMA-ES + noisy CRPs



# Concept - Overview – Motivation [6]



PUF's Category:

PUF's Modeling Attacks with CRPs only:

Weak PUF - small #CRPs:  
RO PUF, SRAM PUF, etc.

Strong PUF – large #CRPs:

Classical ML attacks  
– reliable CRPs:

Advanced ML attacks  
– noisy CRPs:  
CMA-ES + noisy CRPs

**Broken but lightweight:**  
Arbiter PUF/APUF, XOR  
APUF, Feed Forward  
PUF, Lightweight Secure  
PUF, Bistable Ring PUF.

Lightweight,  
Precise Math. Model

Security Proof

Vulnerability

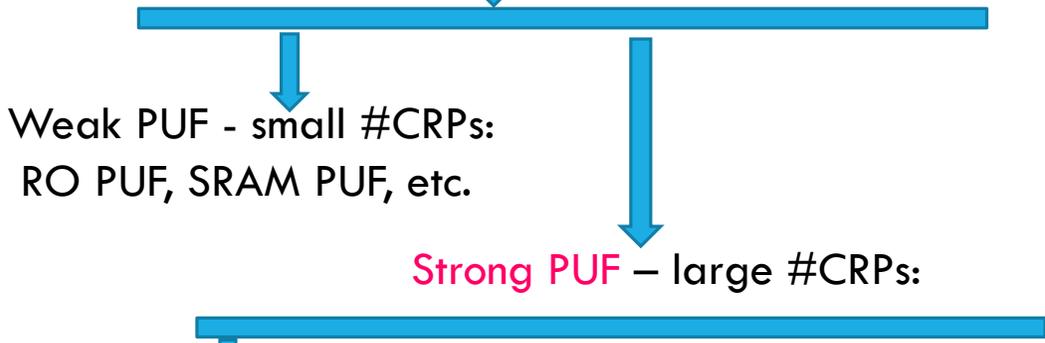
XOR APUF



# Concept - Overview – Motivation [7]



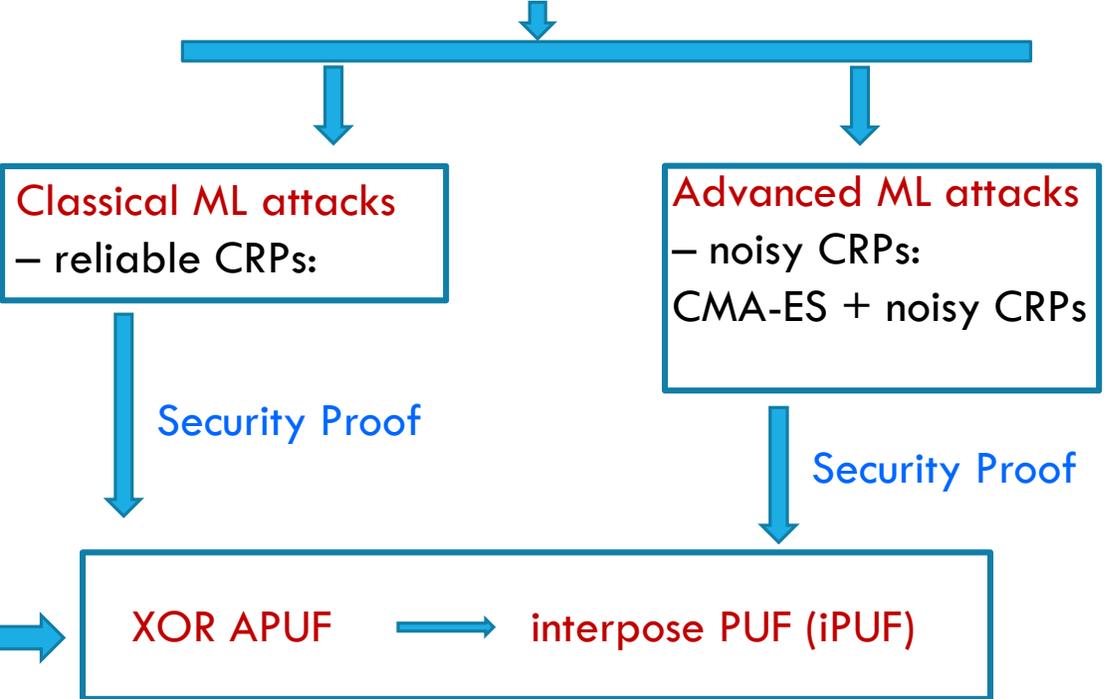
PUF's Category:



**Broken but lightweight:**  
 Arbiter PUF/APUF, XOR APUF, Feed Forward PUF, Lightweight Secure PUF, Bistable Ring PUF.

Lightweight,  
Precise Math. Model

PUF's Modeling Attacks with CRPs only:



# Concept - Overview – Motivation [8]



PUF's Category:

PUF's Modeling Attacks with CRPs only:

Weak PUF - small #CRPs:  
RO PUF, SRAM PUF, etc.

Strong PUF – large #CRPs:

Classical ML attacks  
– reliable CRPs:

Advanced ML attacks  
– noisy CRPs:  
CMA-ES + noisy CRPs

**Broken but lightweight:**  
Arbiter PUF/APUF, XOR APUF, Feed Forward PUF, Lightweight Secure PUF, Bistable Ring PUF.

Lightweight,  
Precise Math. Model

Security Proof

Security Philosophy

Security Proof

XOR APUF

interpose PUF (iPUF)

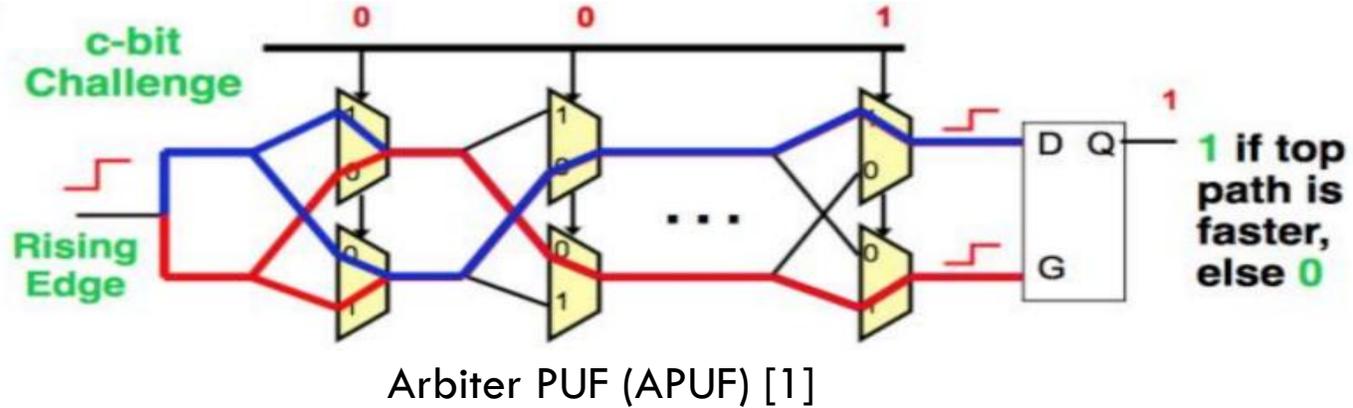
Design Philosophy



## 2. APUF - XOR APUF - iPUF



# APUF, XOR APUF and iPUF [1]

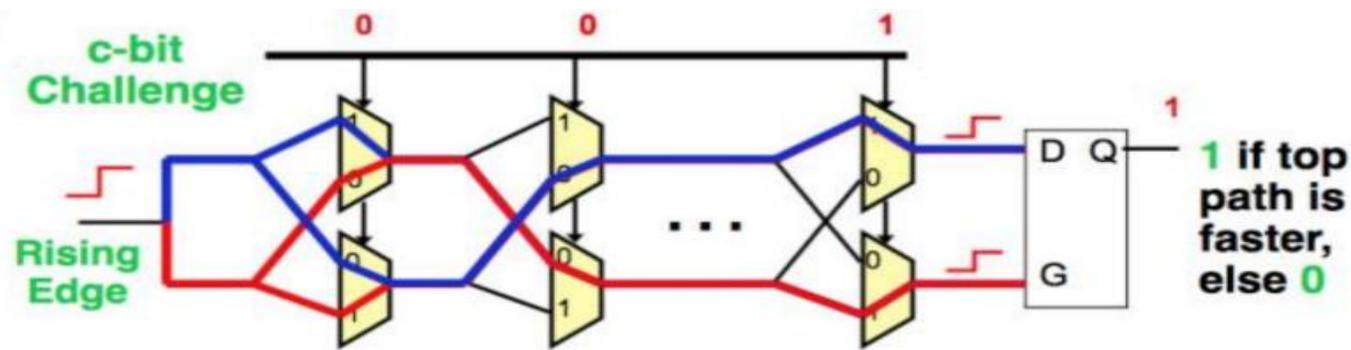


- Extremely lightweight and large number of CRPs i.e,  $2^n$  CRPs
- Environmental noises make the PUF's outputs unreliable sometimes
- Not secure against modeling attacks

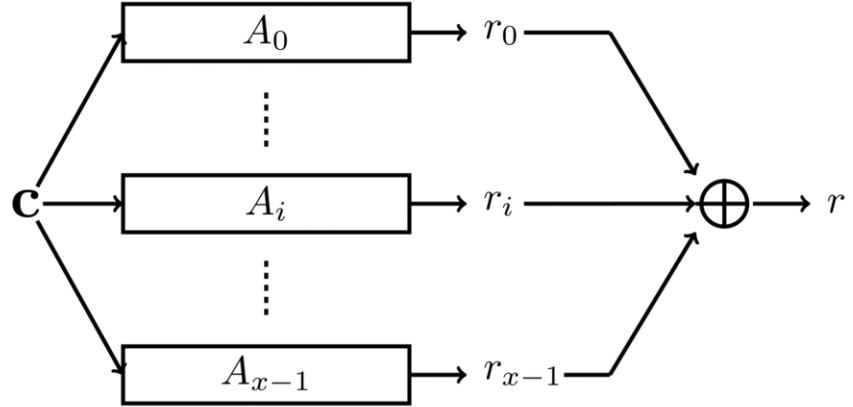


# APUF, XOR APUF and iPUF [2]

Arbiter PUF (APUF)



x-XOR APUF

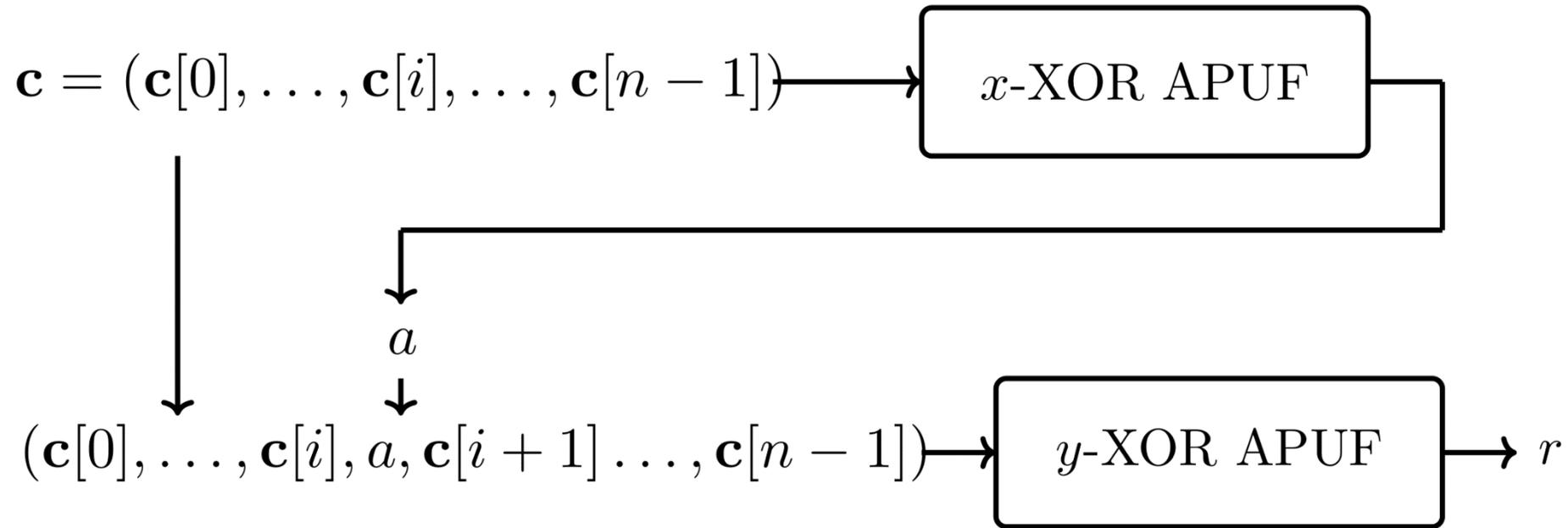


x-XOR PUF.



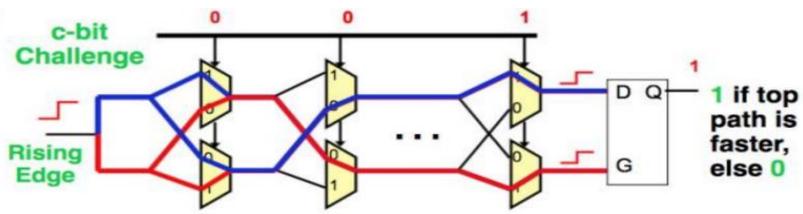
# APUF, XOR APUF and iPUF [3]

The Interpose PUF / iPUF

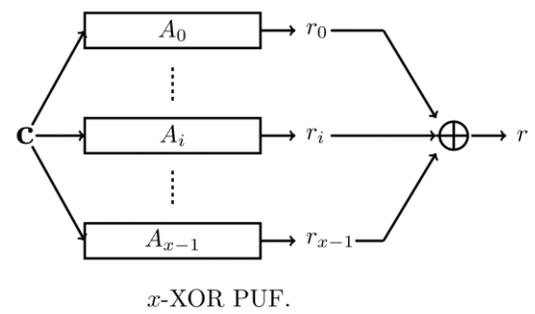


# APUF, XOR APUF and iPUF [4]

Arbiter PUF (APUF)

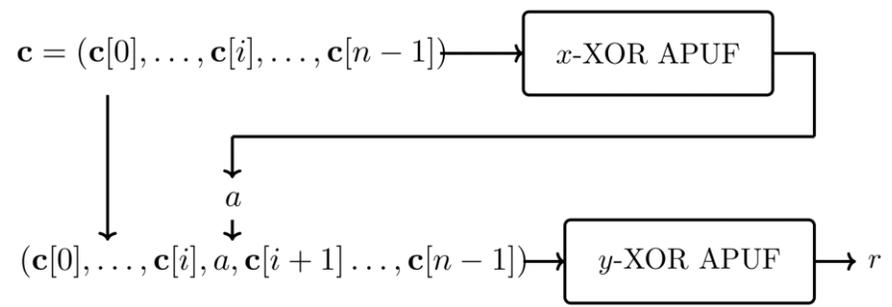


x-XOR Arbiter PUF



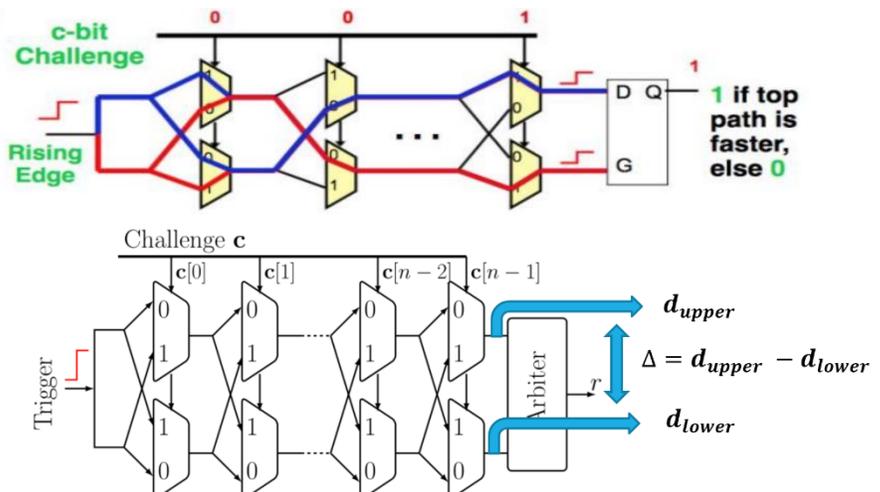
x-XOR PUF.

Interpose PUF (iPUF)



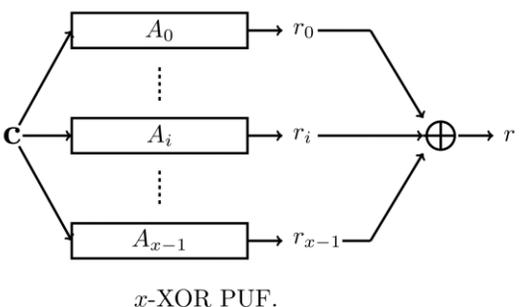
# APUF, XOR APUF and iPUF [5]

Arbiter PUF (APUF)



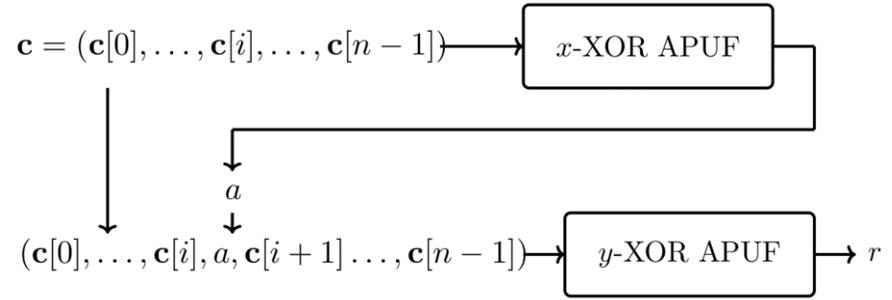
A  $n$ -stage classic Arbiter PUF with challenge  $c \in \{0,1\}^n$ .

x-XOR Arbiter PUF



x-XOR PUF.

Interpose PUF (iPUF)



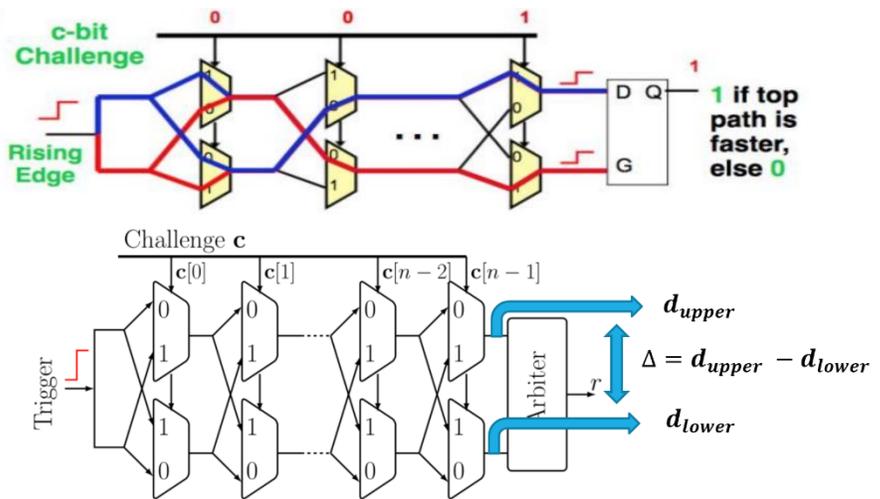
- $\Delta > 0 \rightarrow r = 1$ . Otherwise  $r = 0$
  - $\Delta = d_{upper} - d_{lower} = w \cdot \Phi$
  - $w$  : unique for any APUF instance
  - $\Phi$  is the parity vector
- $$\Phi[i] = \prod_{j=i, \dots, n-1} (1 - c[j]), i = 0, \dots, n-1, \Phi[n] = 1$$



- Precise linear model
- Large CRP space
- Vulnerable to ML attacks

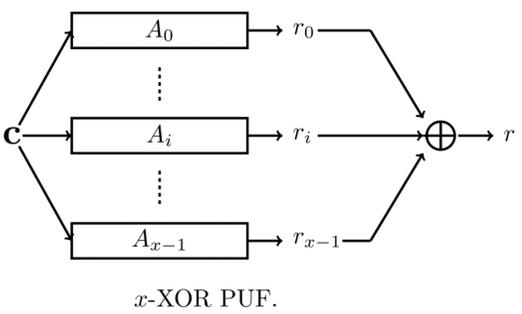
# APUF, XOR APUF and iPUF [6]

Arbiter PUF (APUF)



A  $n$ -stage classic Arbiter PUF with challenge  $\mathbf{c} \in \{0, 1\}^n$ .

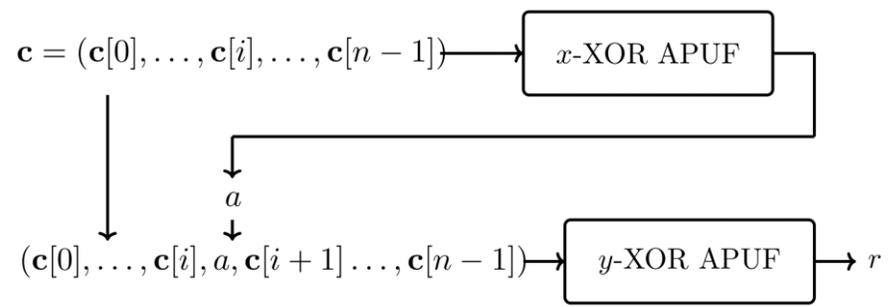
$x$ -XOR Arbiter PUF



$x$ -XOR PUF.

$$r_{\text{XOR APUF}} = \text{sign}\left(\prod_{i=1}^x w_i^T \Phi\right) \quad [2]$$

Interpose PUF (iPUF)



- $\Delta > 0 \rightarrow r = 1$ . Otherwise  $r = 0$
- $\Delta = d_{\text{upper}} - d_{\text{lower}} = \mathbf{w} \cdot \Phi$
- $\mathbf{w}$  : unique for any APUF instance
- $\Phi$  is the parity vector

$$\Phi[i] = \prod_{j=i, \dots, n-1} (1 - c[j]), i = 0, \dots, n - 1, \Phi[n] = 1$$

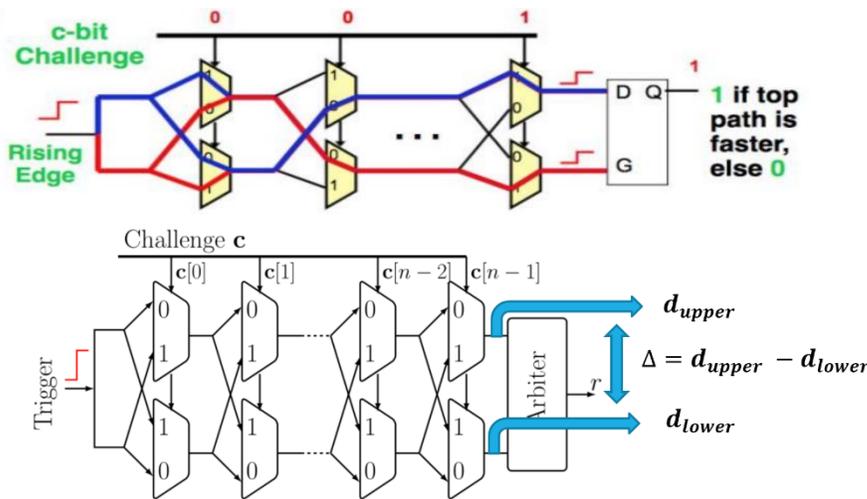
- Precise non-linear model
- Large CRP space
- Secure against classical ML
- Vulnerable to advanced ML



- Precise linear model
- Large CRP space
- Vulnerable to ML attacks

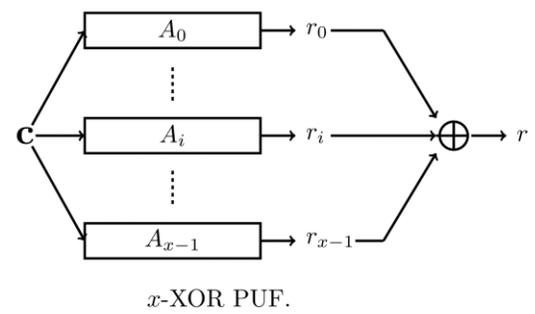
# APUF, XOR APUF and iPUF [7]

Arbiter PUF (APUF)



A  $n$ -stage classic Arbiter PUF with challenge  $\mathbf{c} \in \{0,1\}^n$ .

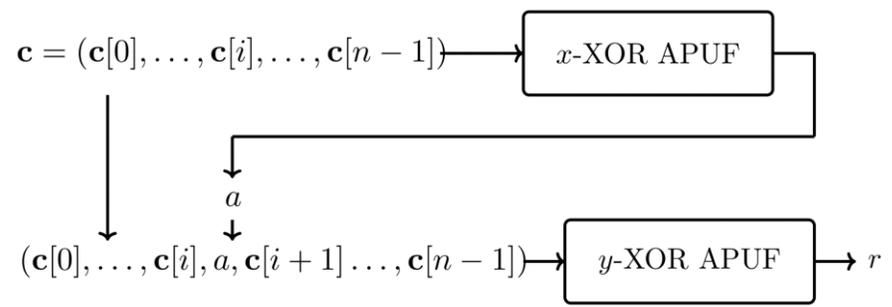
$x$ -XOR Arbiter PUF



$x$ -XOR PUF.

$$r_{\text{XOR APUF}} = \text{sign}\left(\prod_{i=1}^x w_i^T \Phi\right)$$

Interpose PUF (iPUF)



$$(x, y) - \text{IPUF} \approx \left(y + \frac{x}{2}\right) - \text{XOR PUF}$$

if  $a$  is inserted at the middle

- $\Delta > 0 \rightarrow r = 1$ . Otherwise  $r = 0$
- $\Delta = d_{\text{upper}} - d_{\text{lower}} = \mathbf{w} \cdot \Phi$
- $\mathbf{w}$ : unique for any APUF instance
- $\Phi$  is the parity vector

$$\Phi[i] = \prod_{j=i, \dots, n-1} (1 - c[j]), i = 0, \dots, n-1, \Phi[n] = 1$$

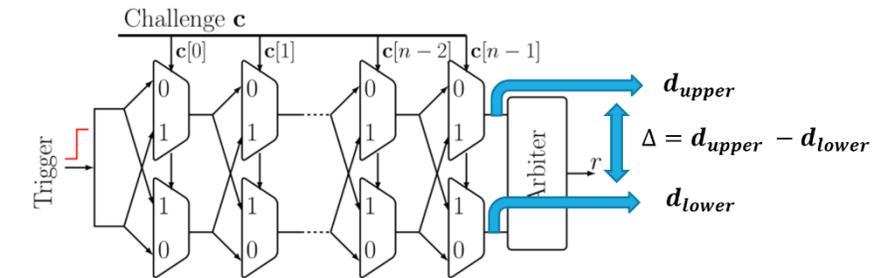
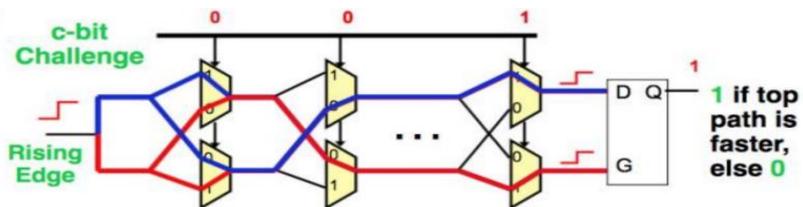
- Precise non-linear model
- Large CRP space
- Secure against classical ML
- Vulnerable to advanced ML



- Precise linear model
- Large CRP space
- Vulnerable to ML attacks

# APUF, XOR APUF and iPUF [8]

Arbiter PUF (APUF)

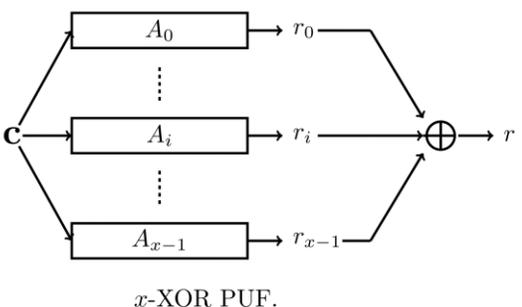


A n-stage classic Arbiter PUF with challenge  $c \in \{0,1\}^n$ .

- $\Delta > 0 \rightarrow r = 1$ . Otherwise  $r = 0$
- $\Delta = d_{upper} - d_{lower} = w \cdot \Phi$
- $w$ : unique for any APUF instance
- $\Phi$  is the parity vector

$$\Phi[i] = \prod_{j=i, \dots, n-1} (1 - c[j]), i = 0, \dots, n - 1, \Phi[n] = 1$$

x-XOR Arbiter PUF

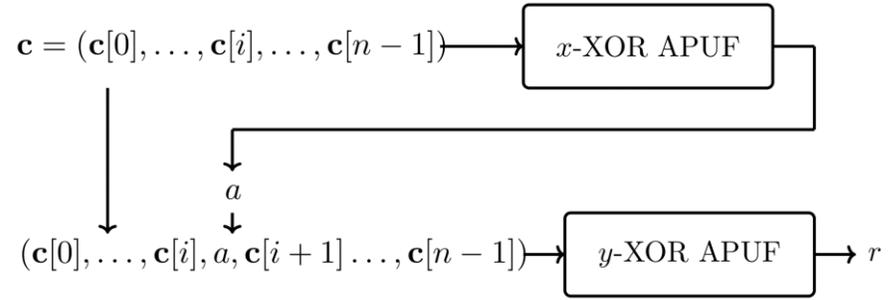


x-XOR PUF.

$$r_{XOR APUF} = sign\left(\prod_{i=1}^x w_i^T \Phi\right)$$

- Precise non-linear model
- Large CRP space
- Secure against classical ML
- Vulnerable to advanced ML

Interpose PUF (iPUF)



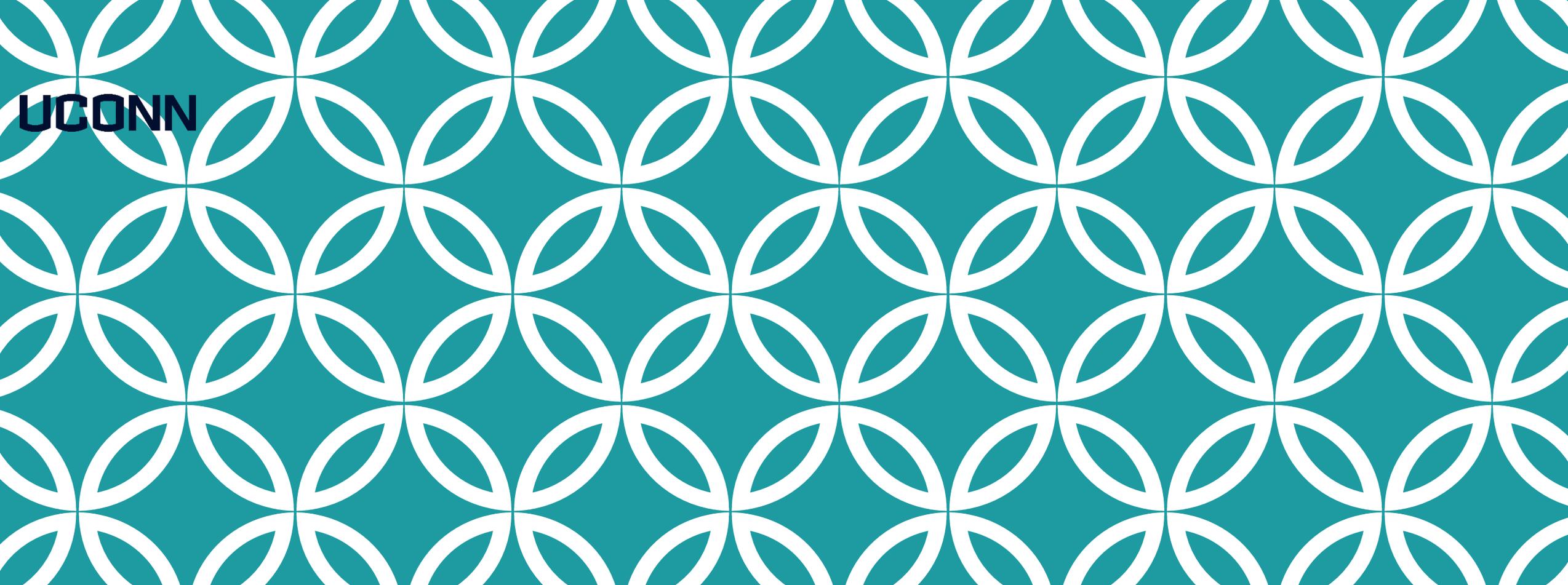
$$(x, y) - IPUF \approx \left(y + \frac{x}{2}\right) - XOR PUF$$

if a is inserted at the middle

- Precise non-linear model
- Large CRP
- Secure both classical ML and advanced ML



- Precise linear model
- Large CRP space
- Vulnerable to ML attacks

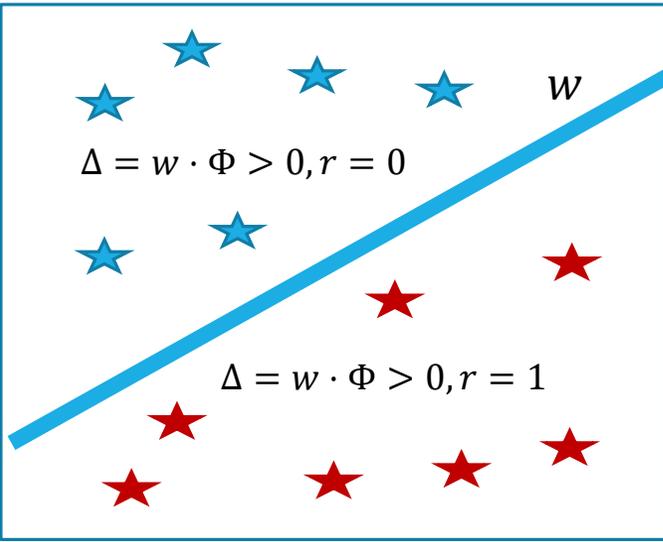
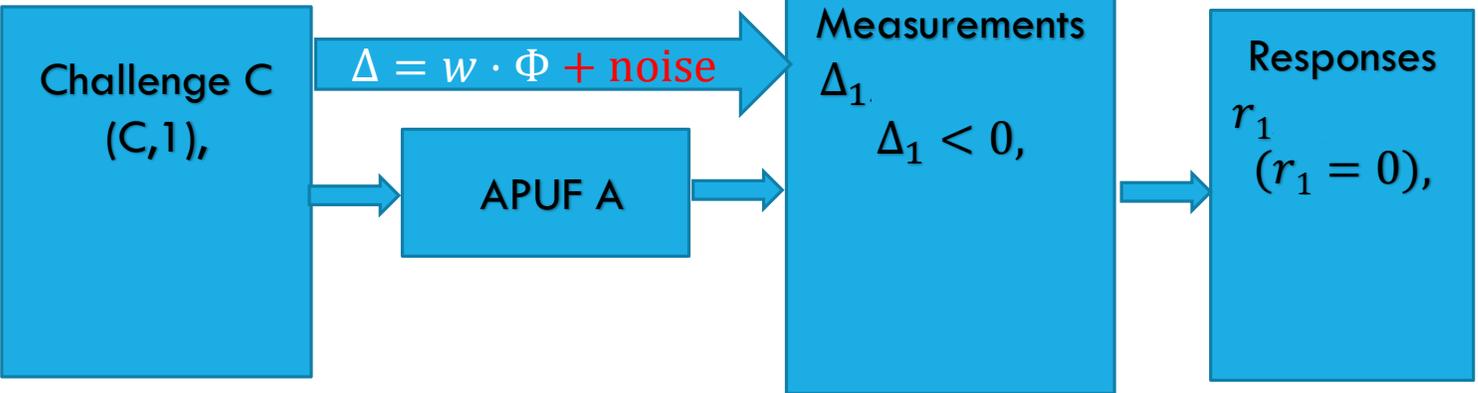


**UConn**

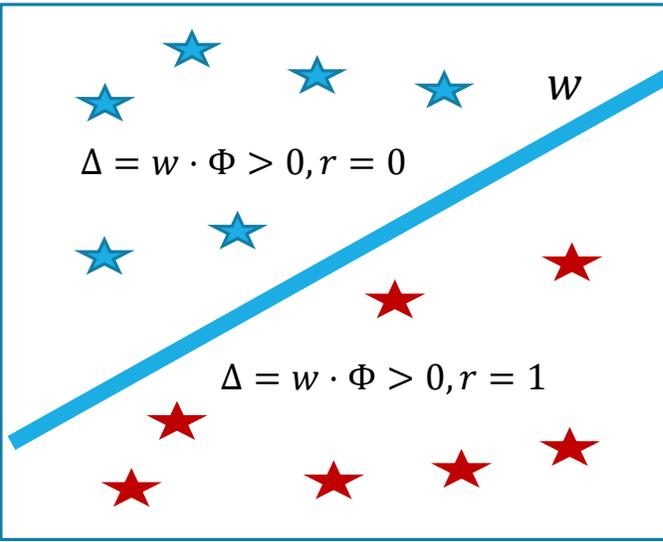
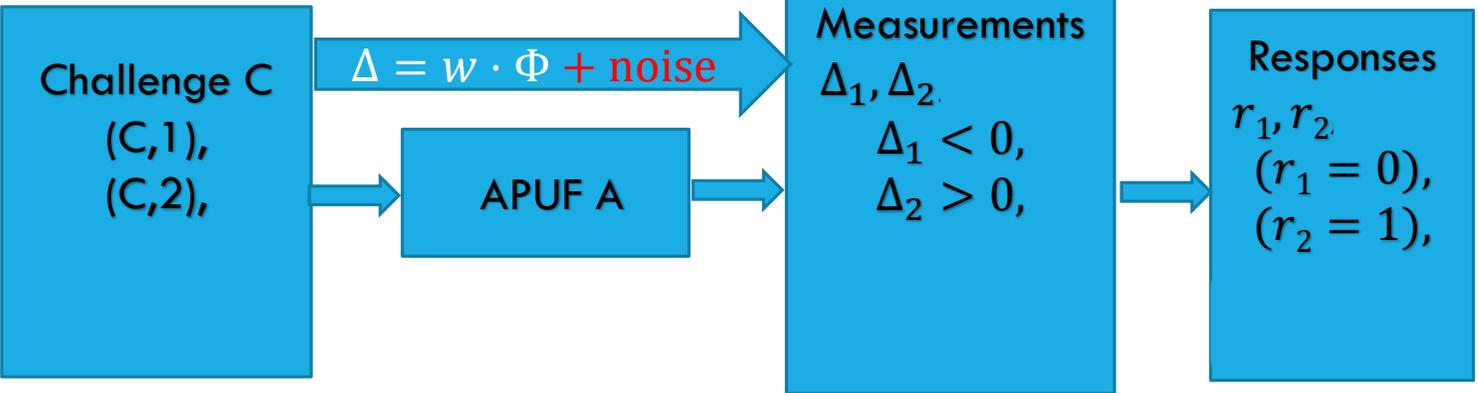
### **3. Short-term Reliability** |



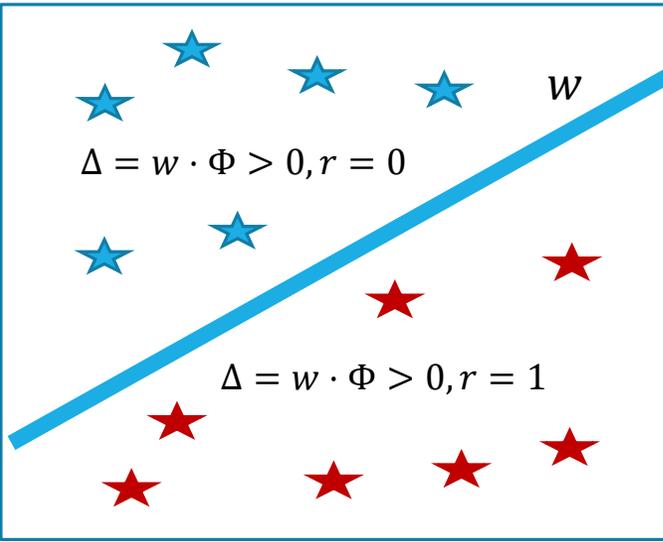
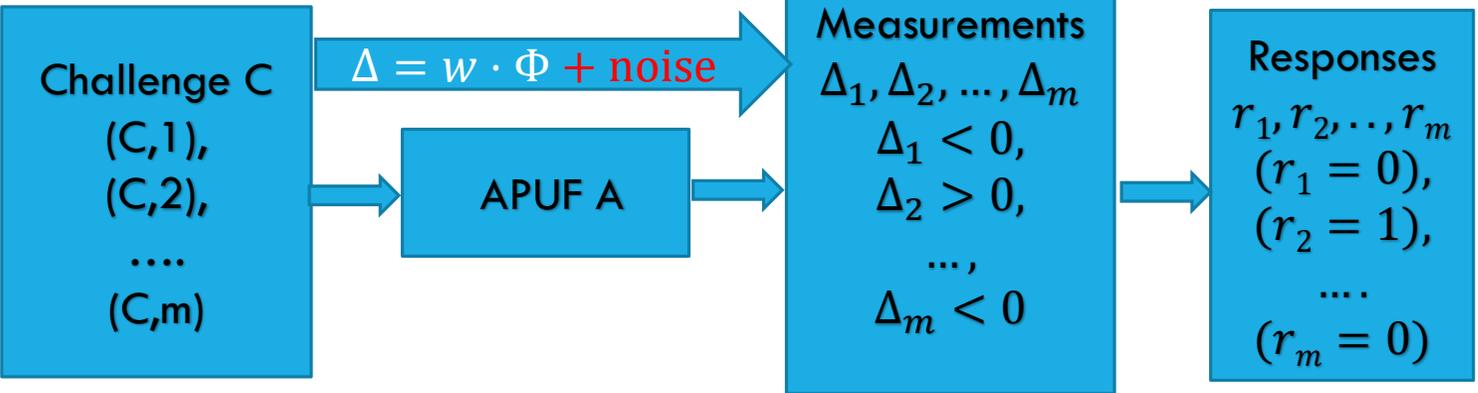
# Arbiter: Repeatability – short-term Reliability [1]



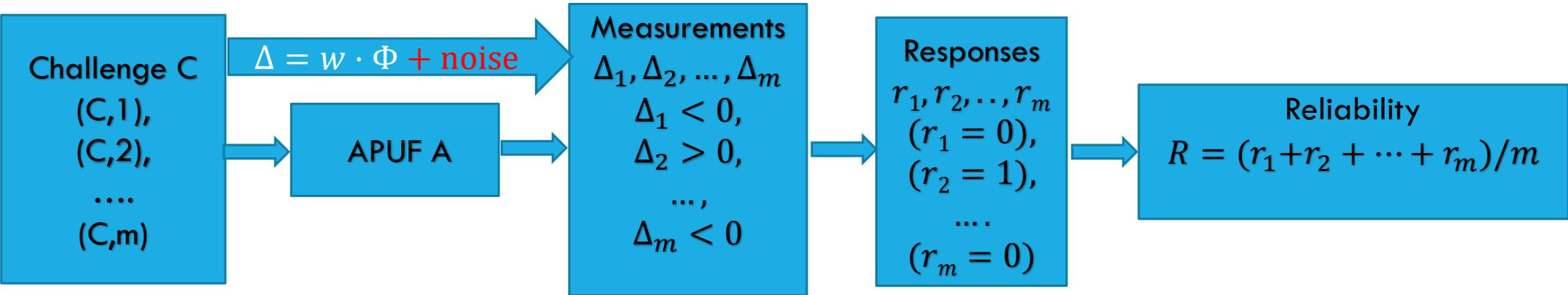
# Arbiter: Repeatability – short-term Reliability [2]



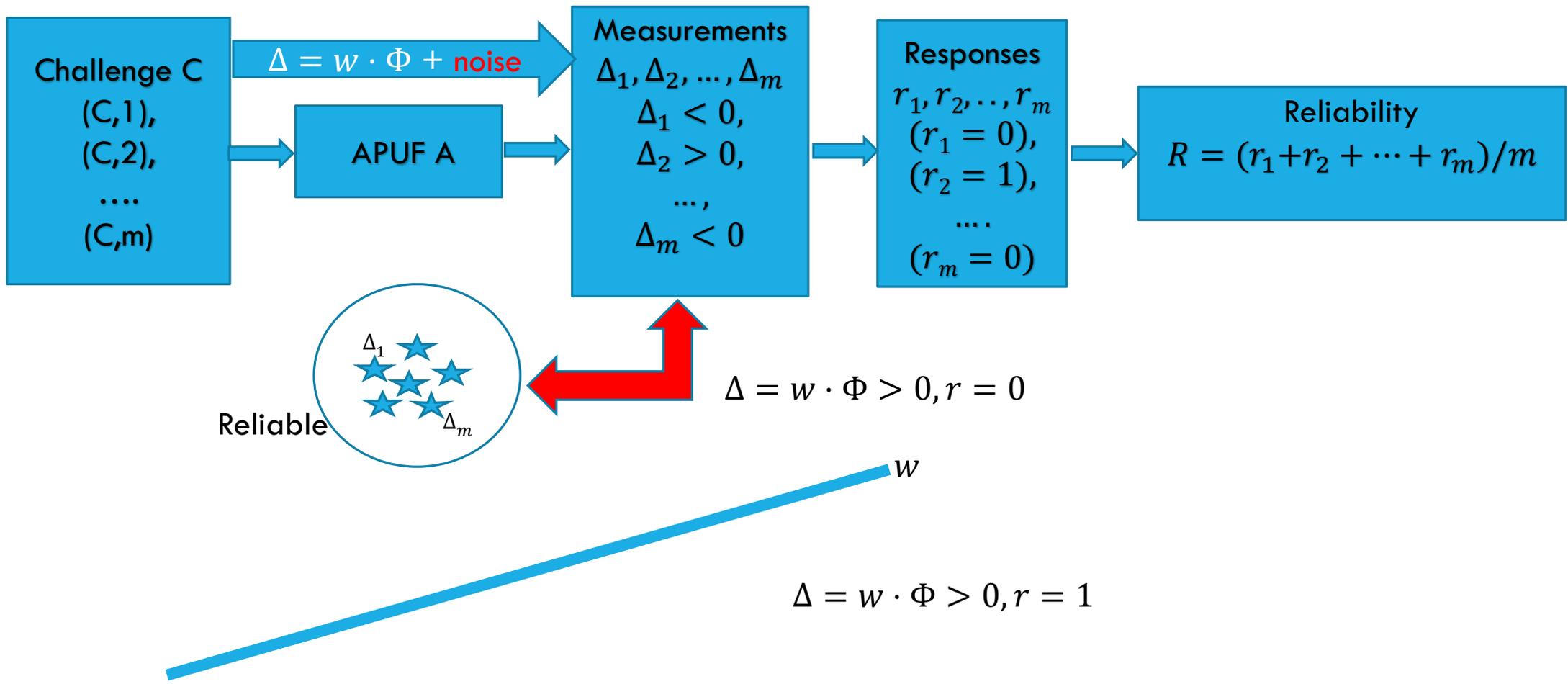
# Arbiter: Repeatability – short-term Reliability [3]



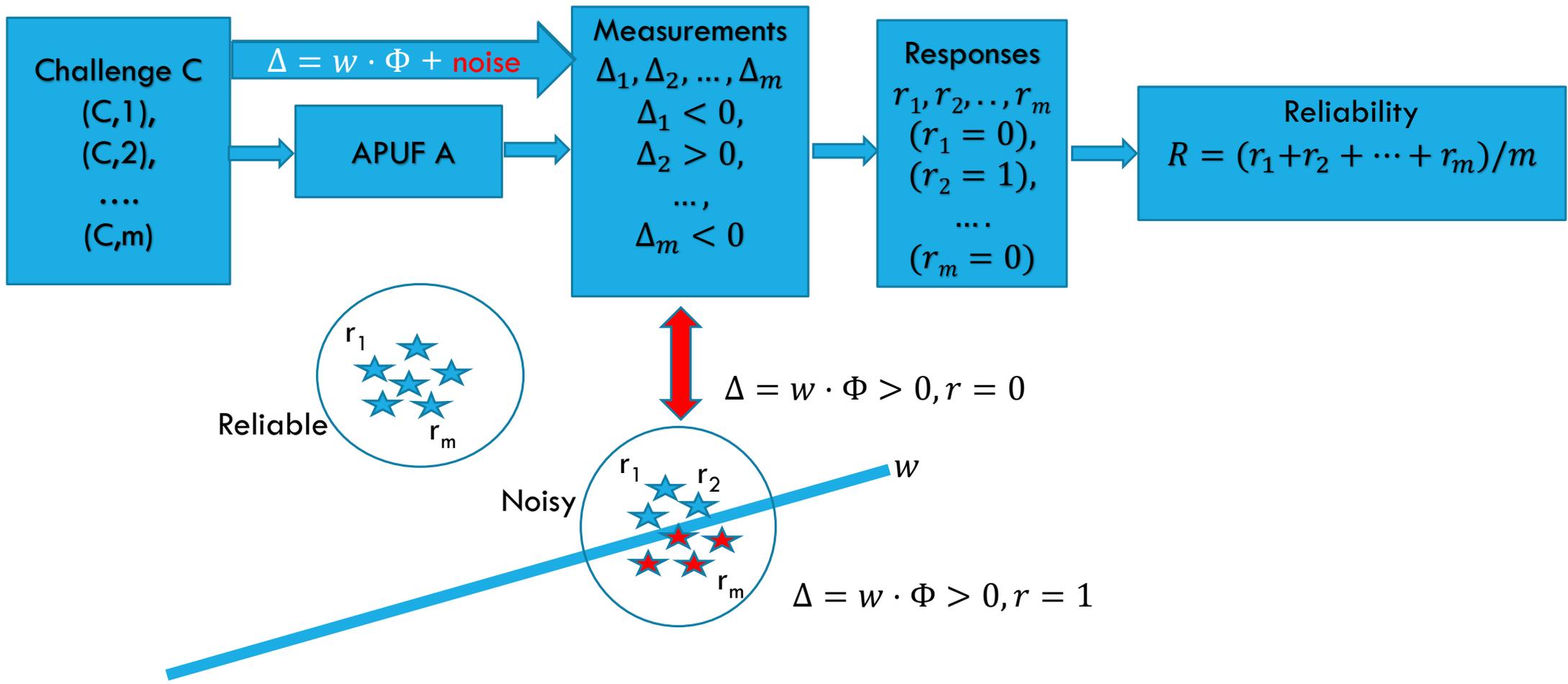
# Arbiter: Repeatability – short-term Reliability [4]



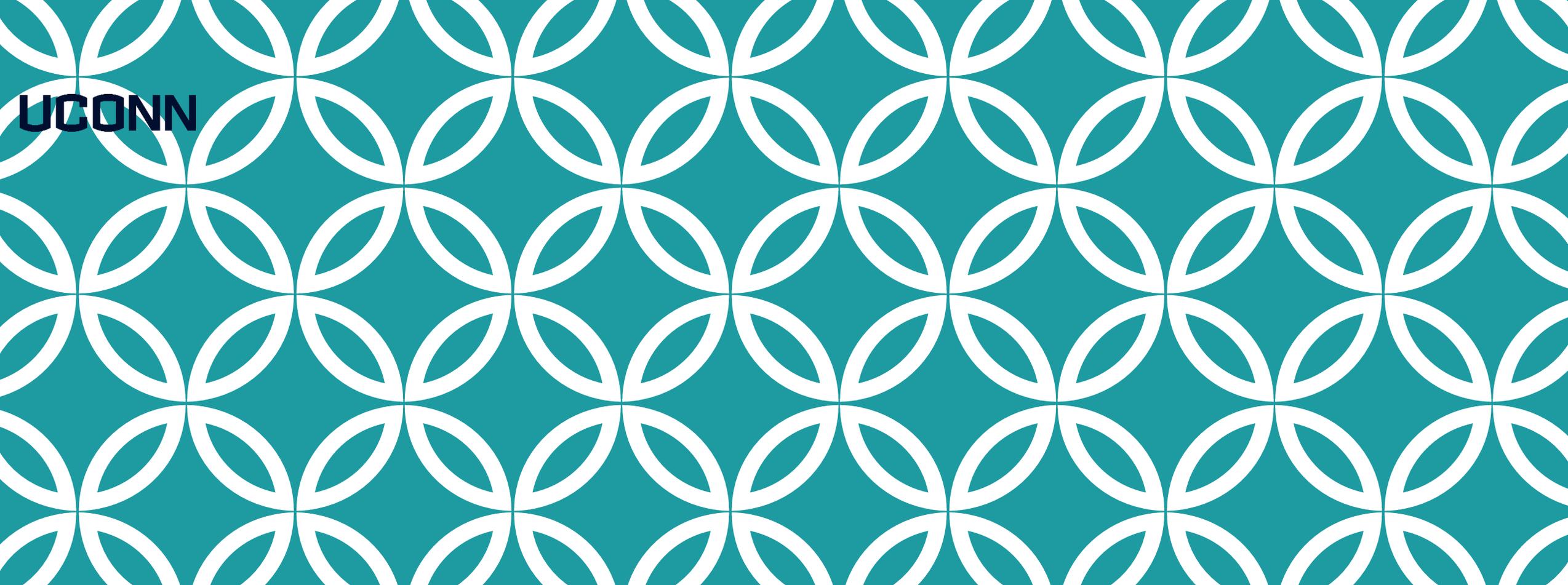
# Arbiter: Repeatability – short-term Reliability [5]



# Arbiter: Repeatability – short-term Reliability [6]





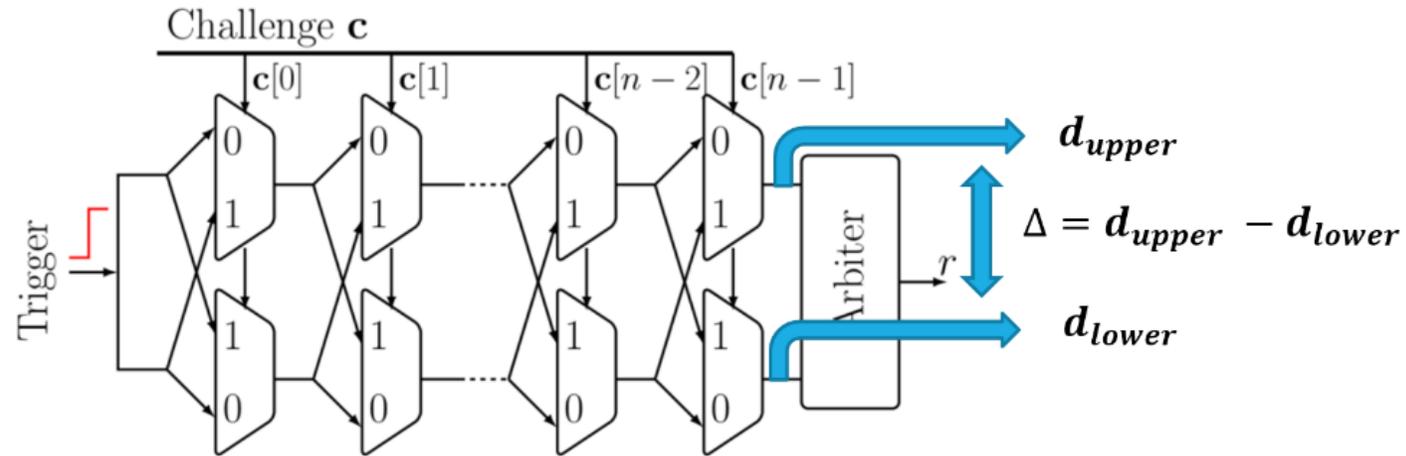


**UConn**

# 4. Reliability based Modeling Attacks



# APUF



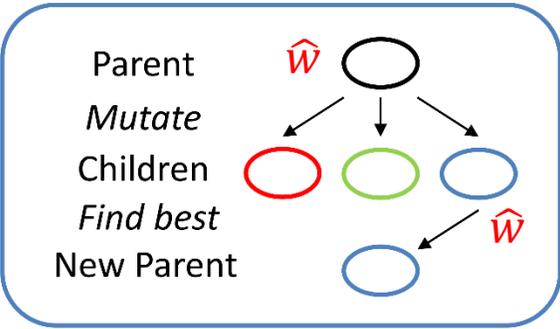
A  $n$ -stage classic Arbiter PUF with challenge  $\mathbf{c} \in \{0, 1\}^n$ .

- $\Delta > 0 \rightarrow r = 1$ . Otherwise  $r = 0$
- $\Delta = d_{upper} - d_{lower} = \mathbf{w} \cdot \Phi$



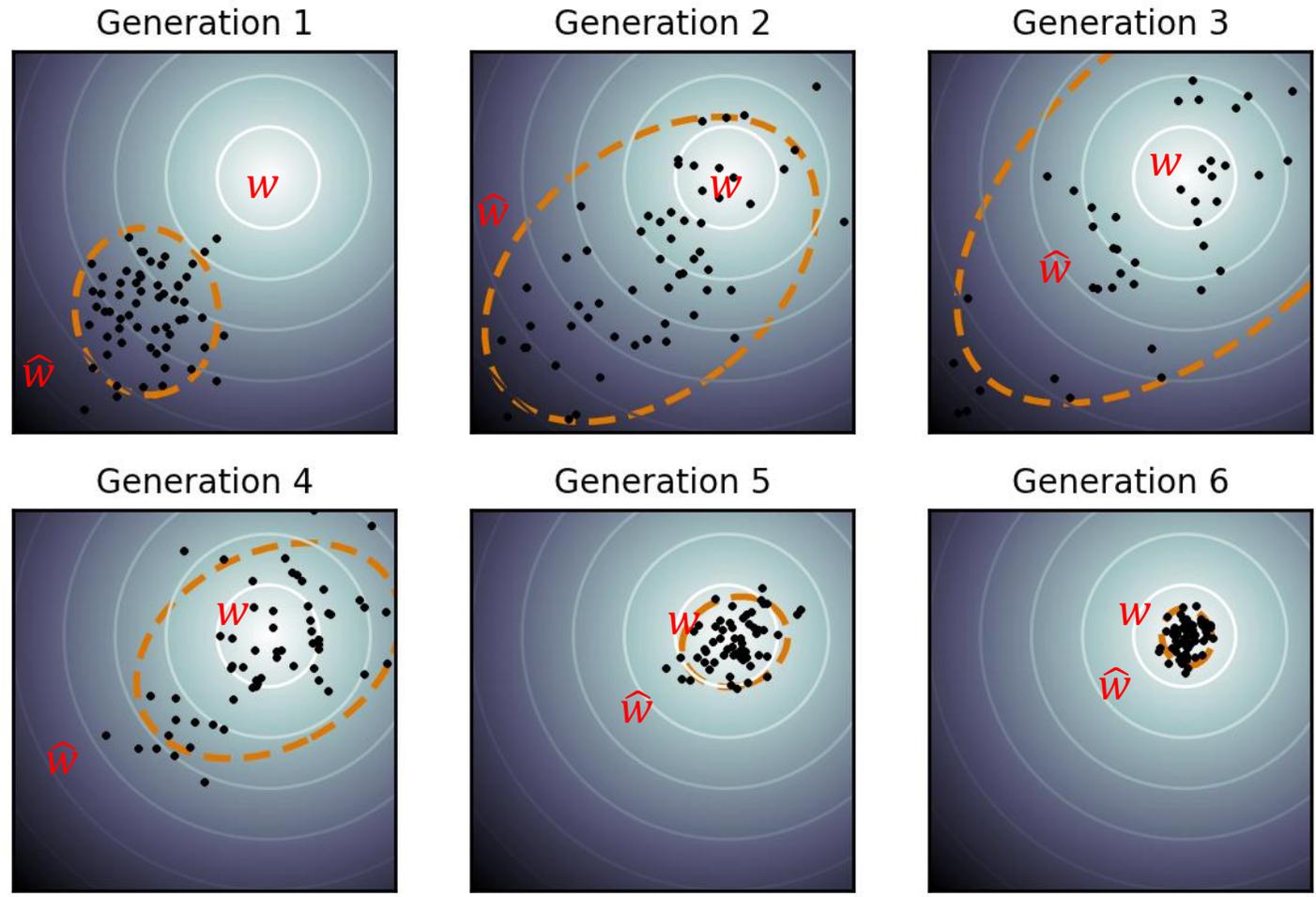
$w$ : target

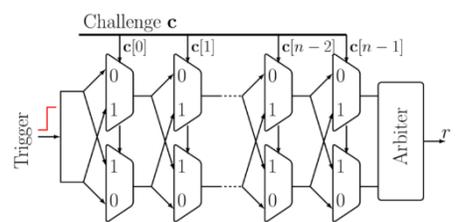
$\hat{w}$ : estimator or model



[4]

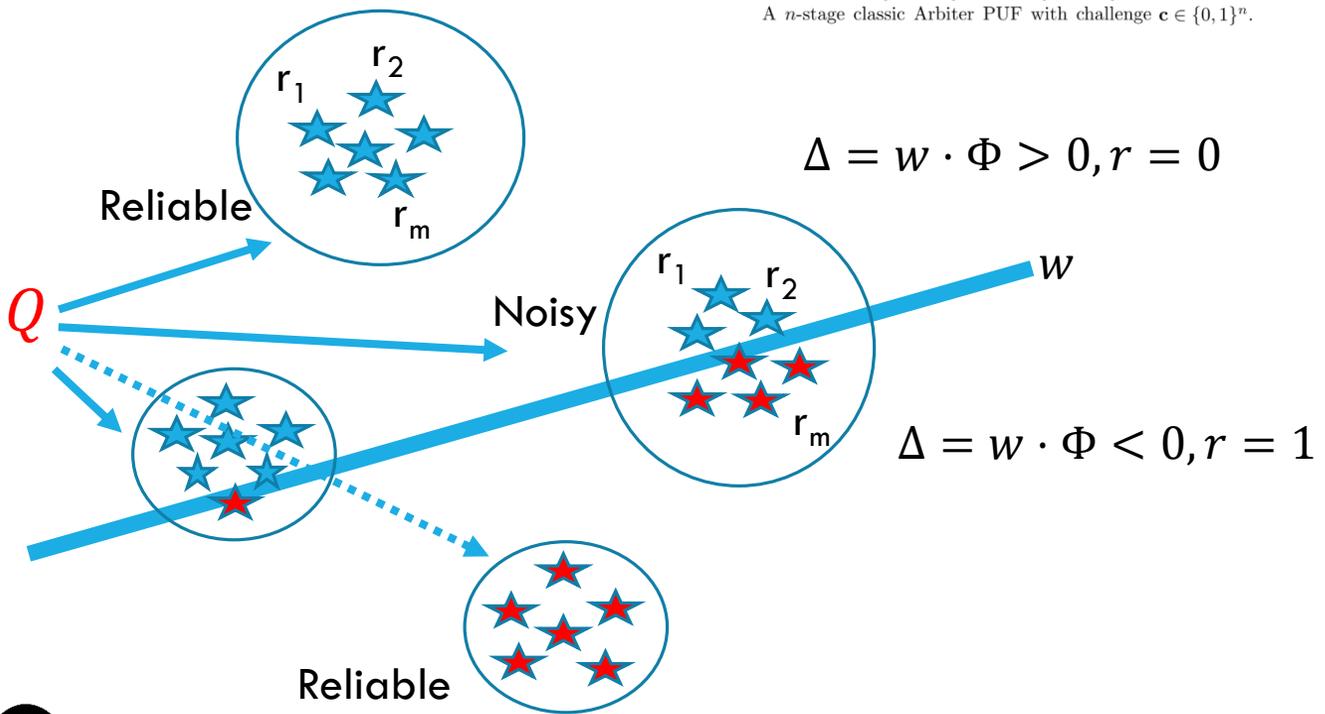
$$\hat{w} = w$$





A  $n$ -stage classic Arbiter PUF with challenge  $c \in \{0, 1\}^n$ .

$Q$ : set of CRPs ,  
 $w$ : APUF



Iteration 1

Target

$Q, w$



Model

$Q, \epsilon_1, \hat{w}_1$

$Q \rightarrow \text{challenge } c \rightarrow \Phi(c) \rightarrow \Delta = \hat{w}_1 \cdot \Phi(c)$

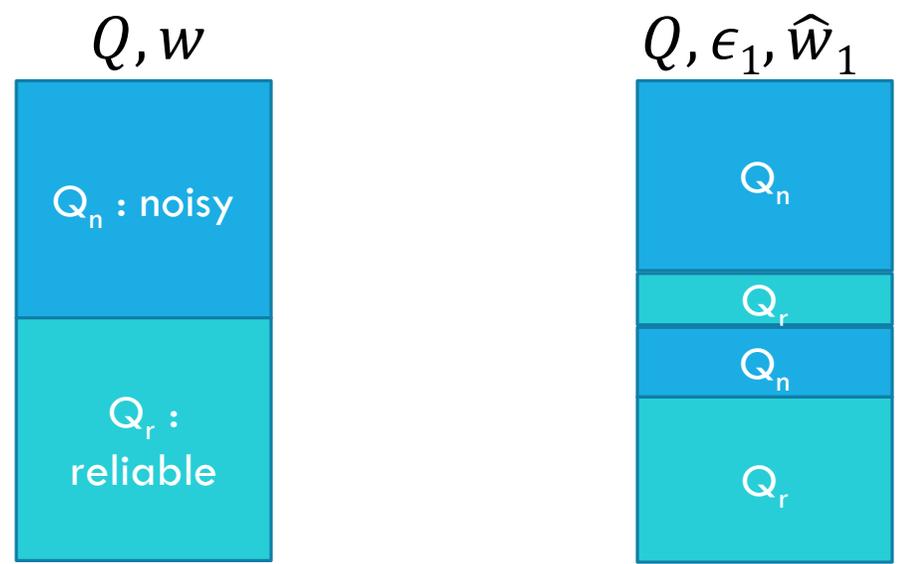
$|\Delta| \leq \epsilon_1 \rightarrow \text{challenge } c \text{ is noisy}$



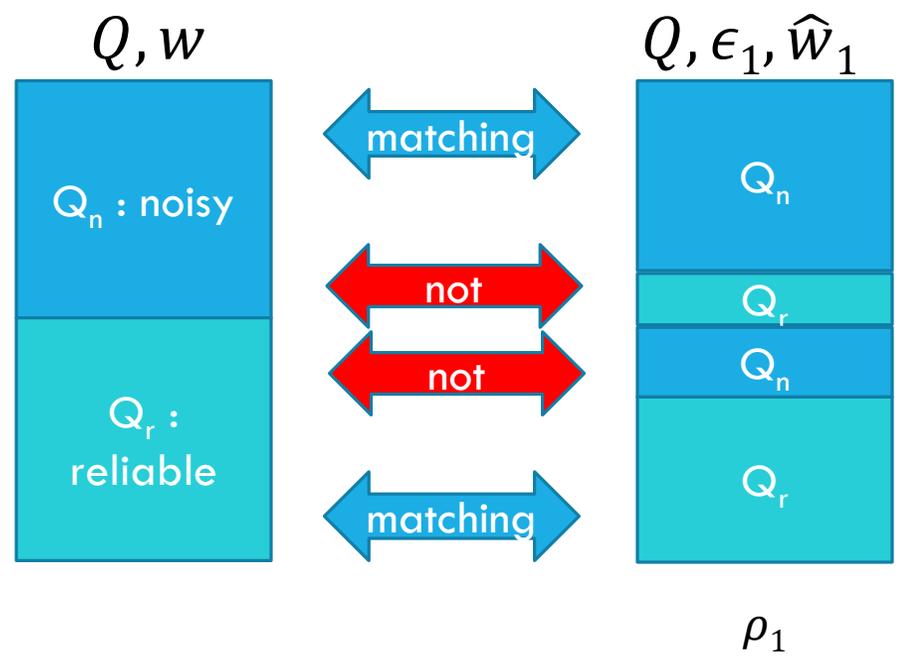
$|\Delta| > \epsilon_1 \rightarrow \text{challenge } c \text{ is reliable}$



Iteration 1



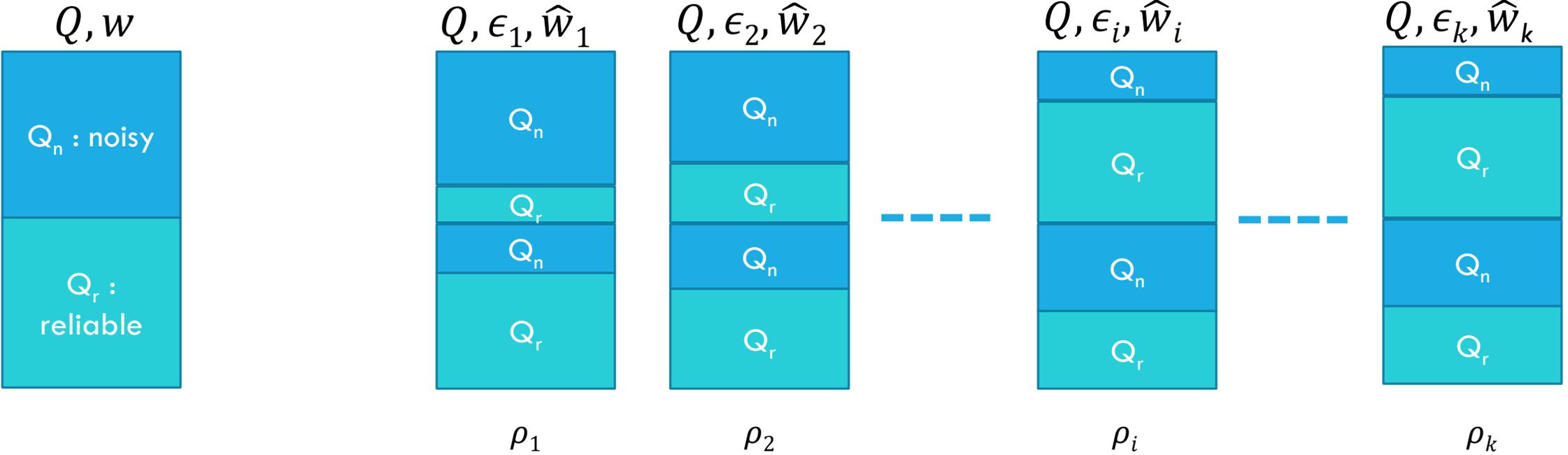
Iteration 1

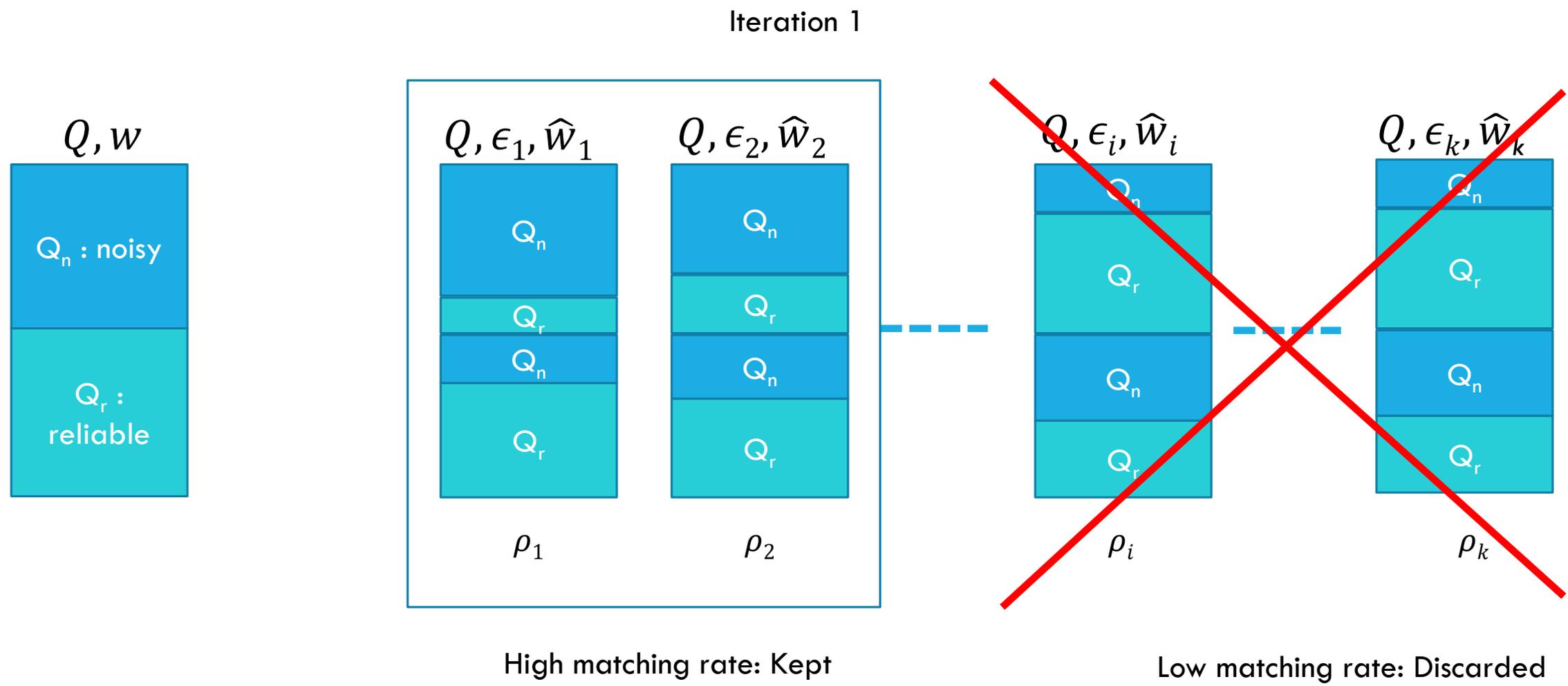


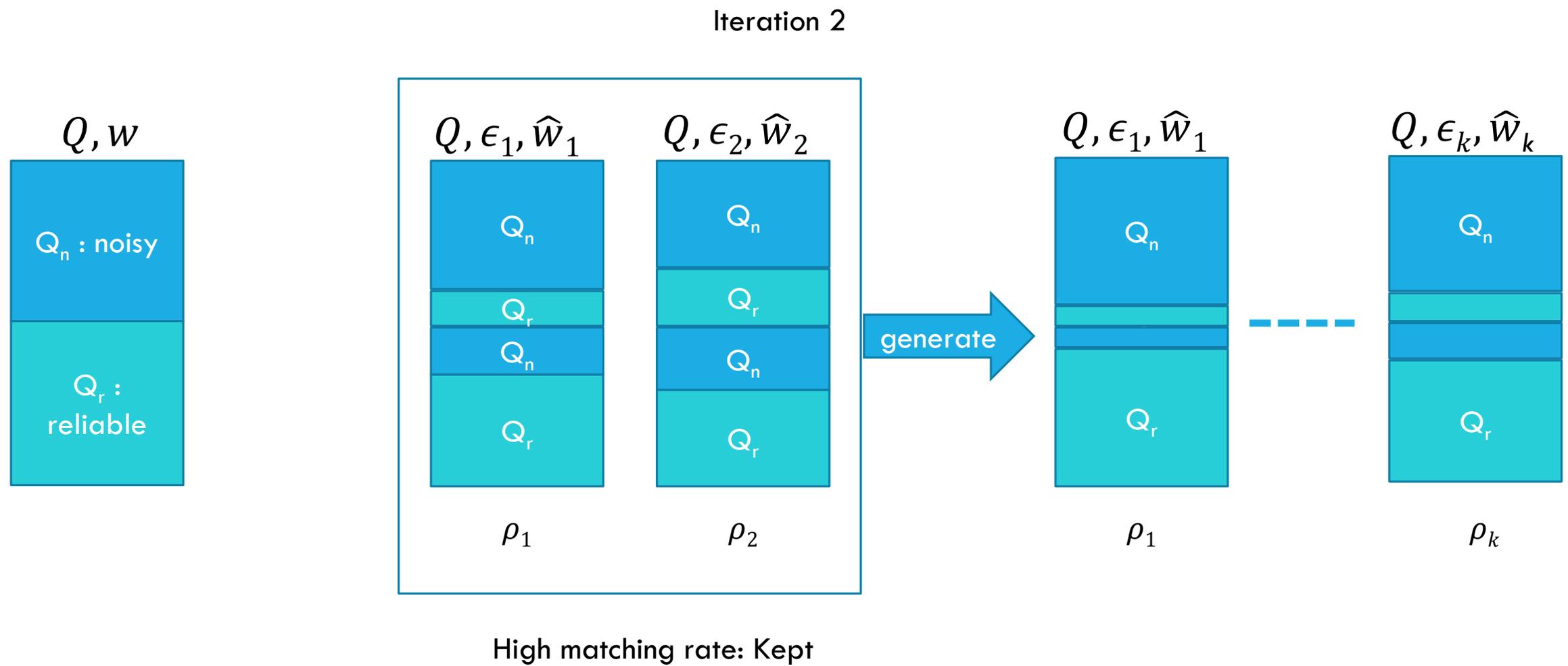
Compute the matching rate  $\rho_1$  between  $[Q, w]$  and  $[Q, \epsilon_1, \hat{w}_1]$

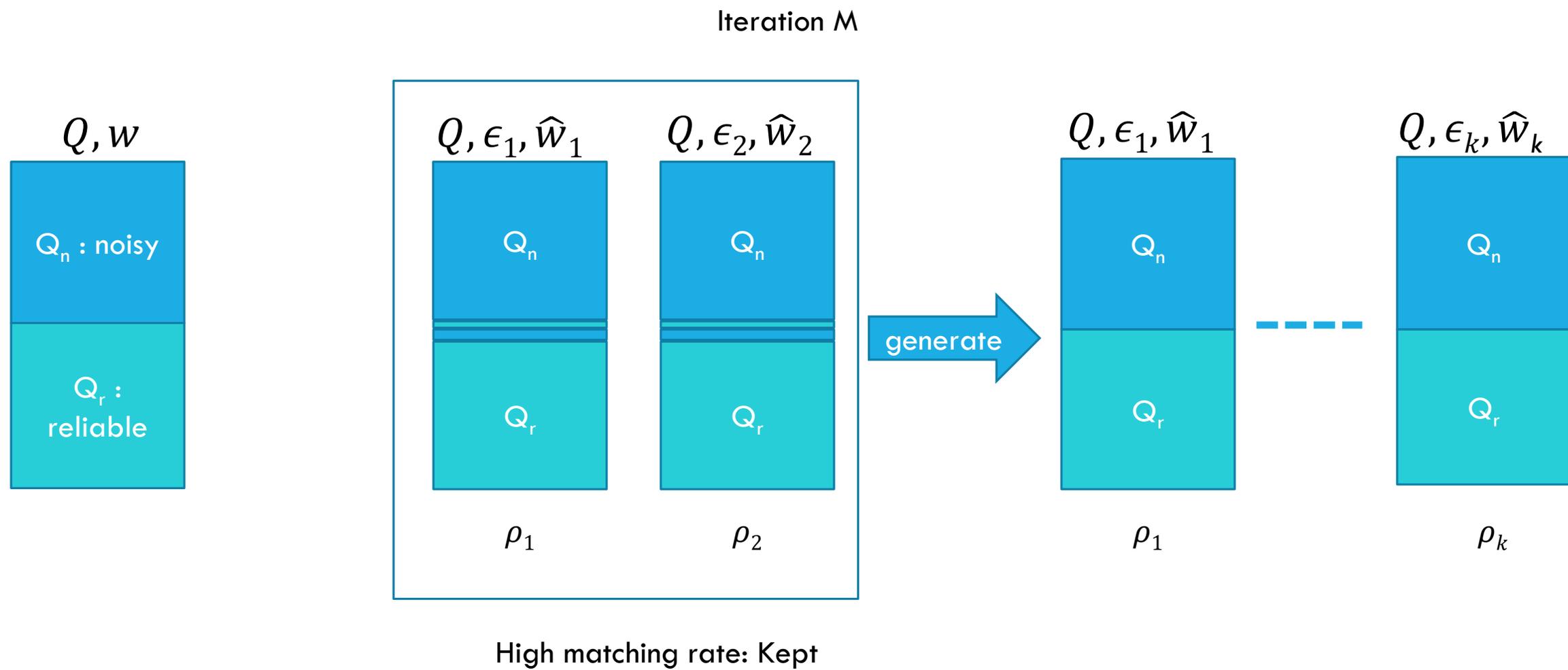


Iteration 1

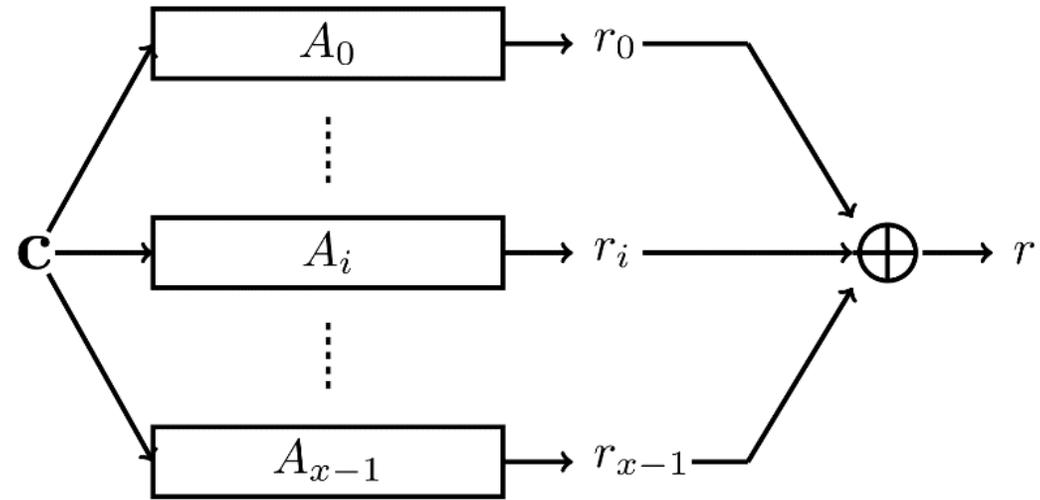








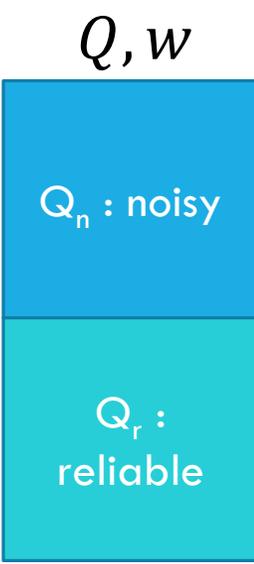
# $x$ -XOR APUF



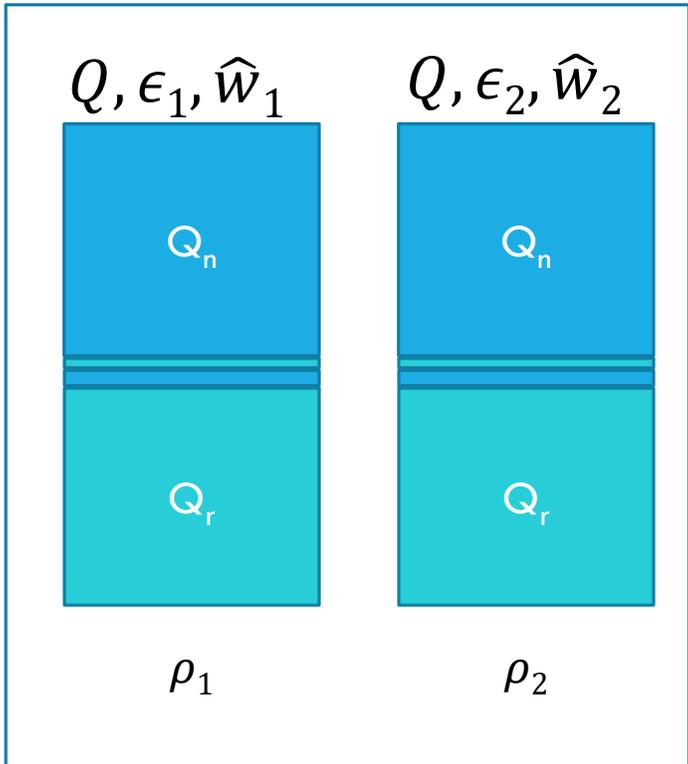
$x$ -XOR PUF.



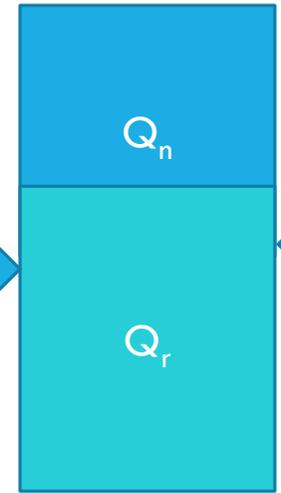
XOR APUF  $\rightarrow Q$



Iteration M



$\hat{W}_1, \dots, \hat{W}_k$  are STILL models of APUF

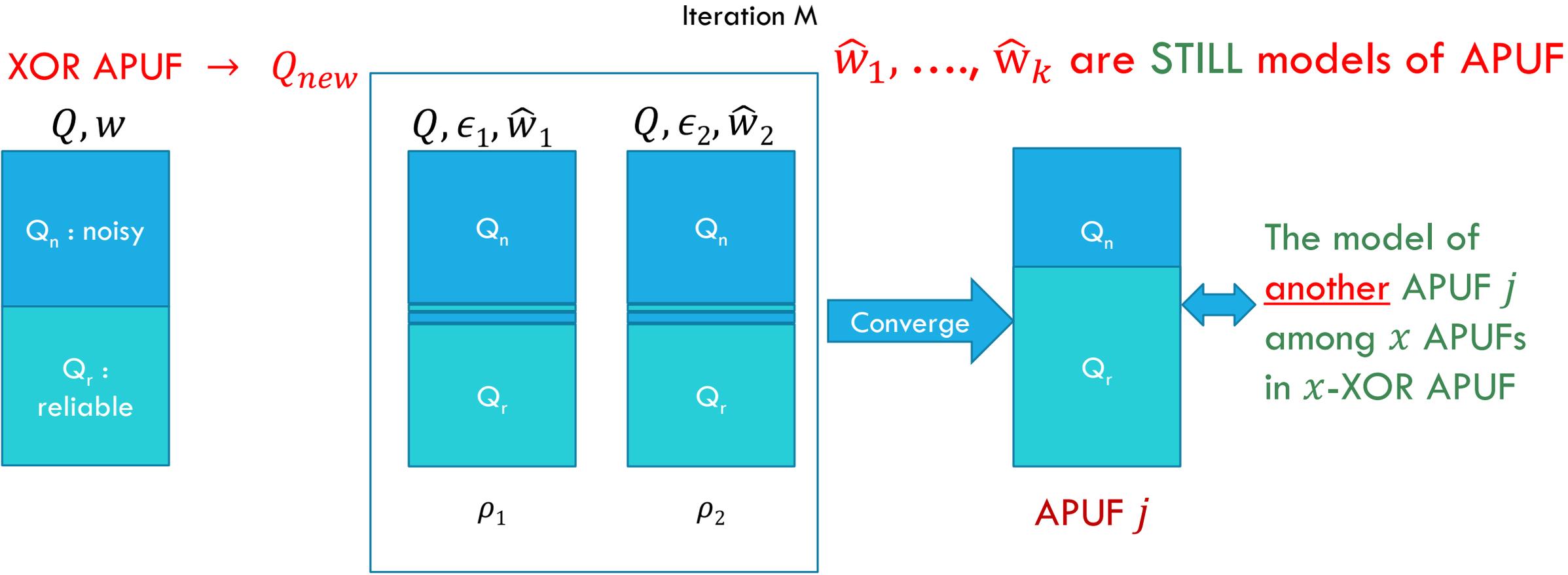


The model of one APUF  $i$  among  $x$  APUFs in  $x$ -XOR APUF

High matching rate: Kept

CMA-ES attack on XOR APUF





- Question 1: How does the attack on XOR PUF work?
- Question 2: How can we make the attack on XOR PUF fail?

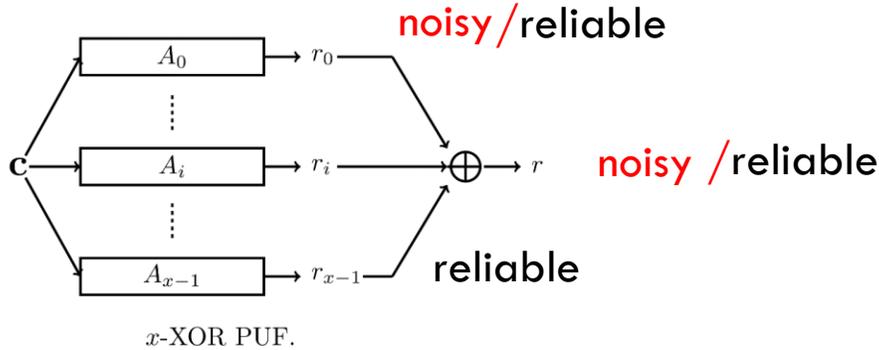


# Question 1: How does the attack on XOR PUF work?

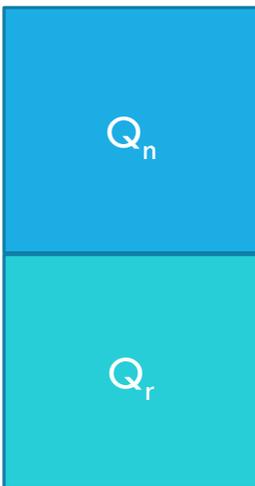
---



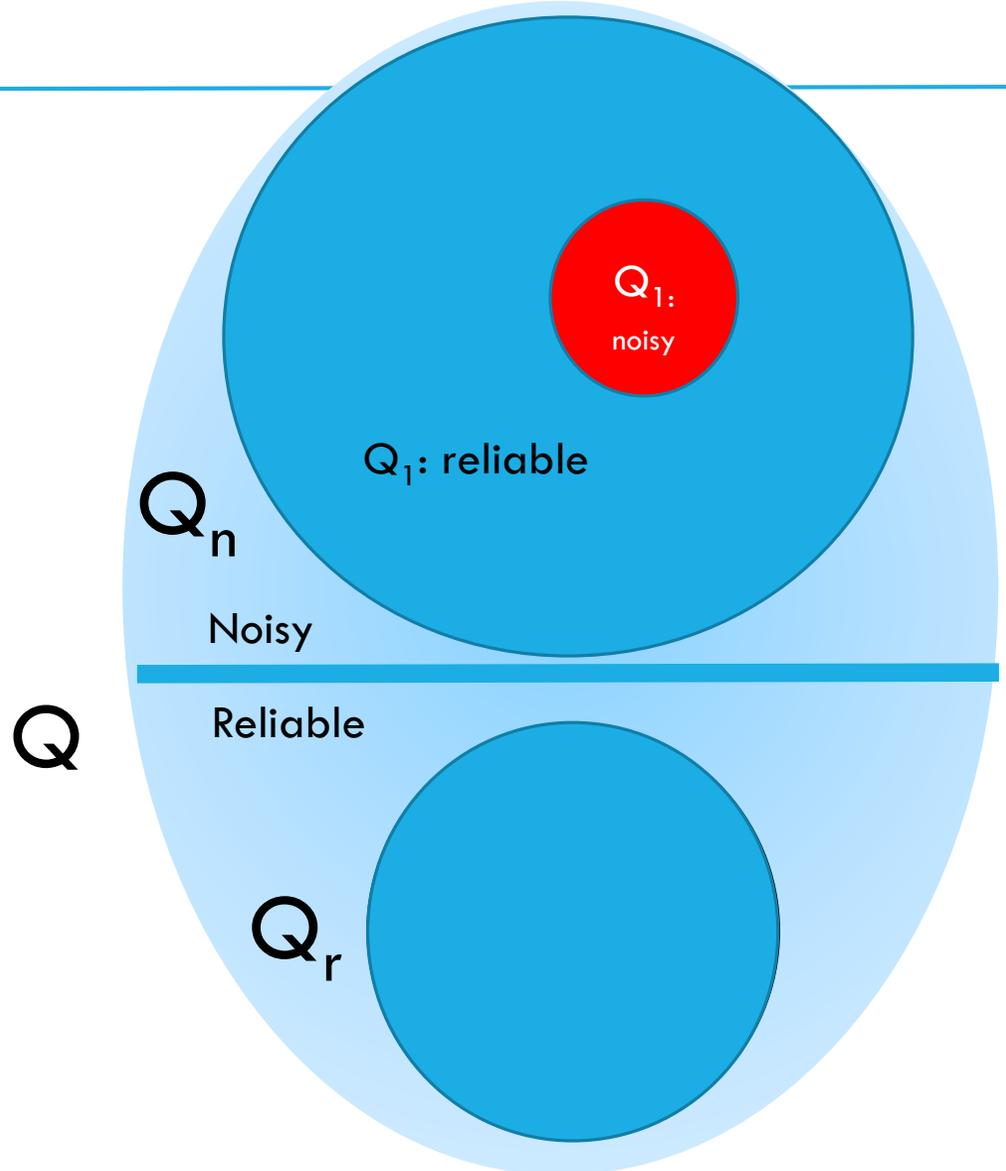
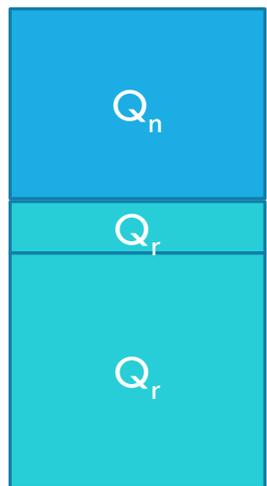
# The noisy and reliable challenges in XOR PUF



XOR APUF

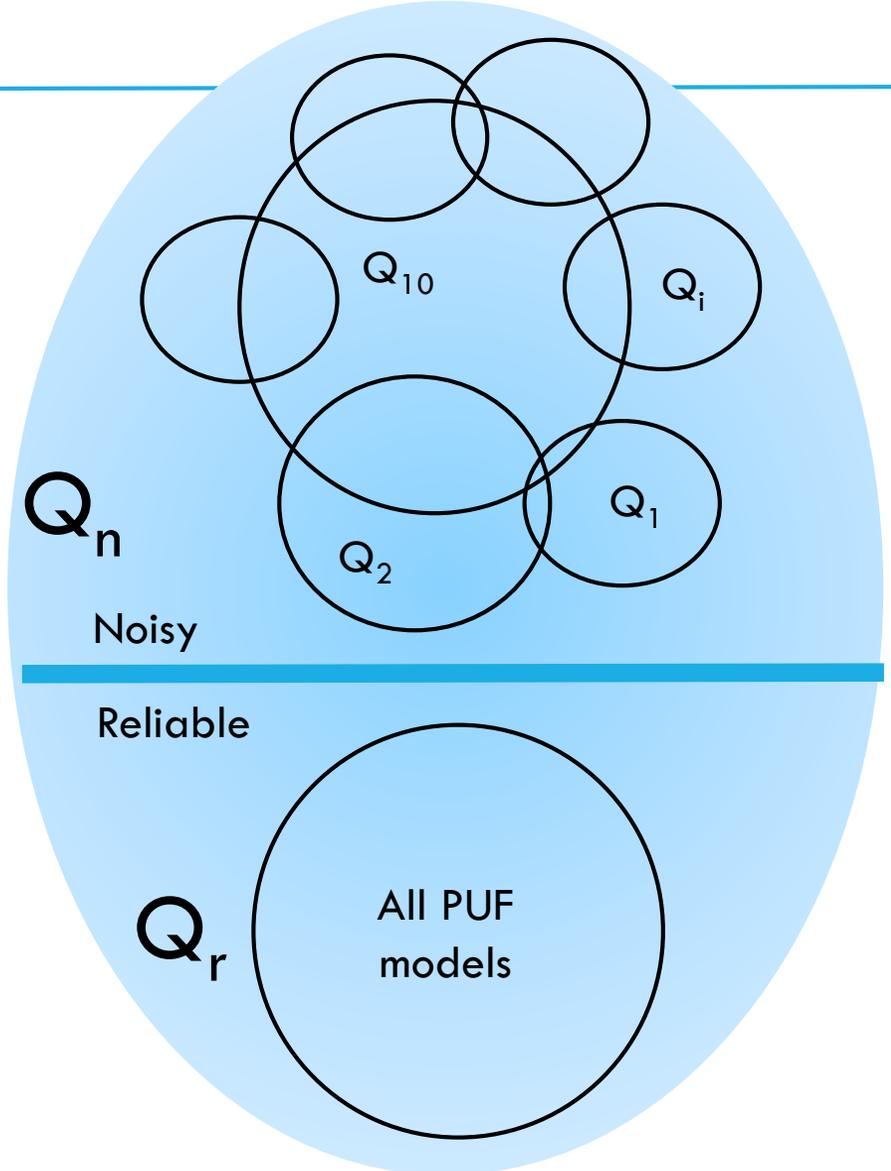
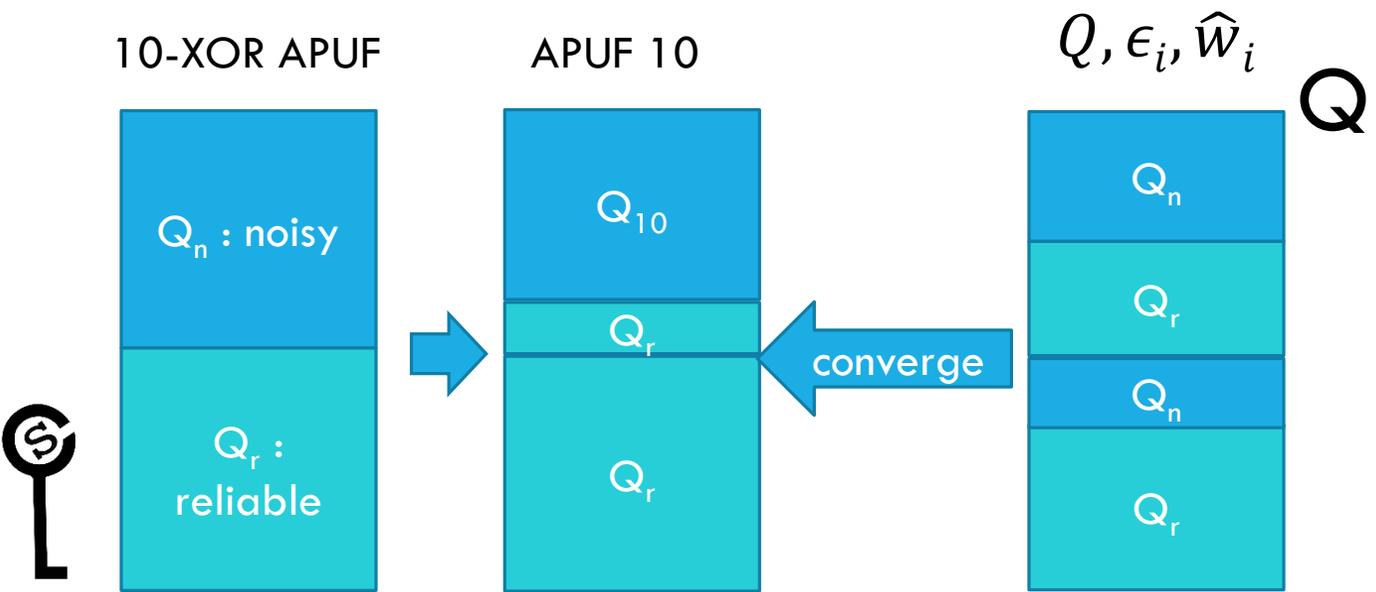


APUF 1



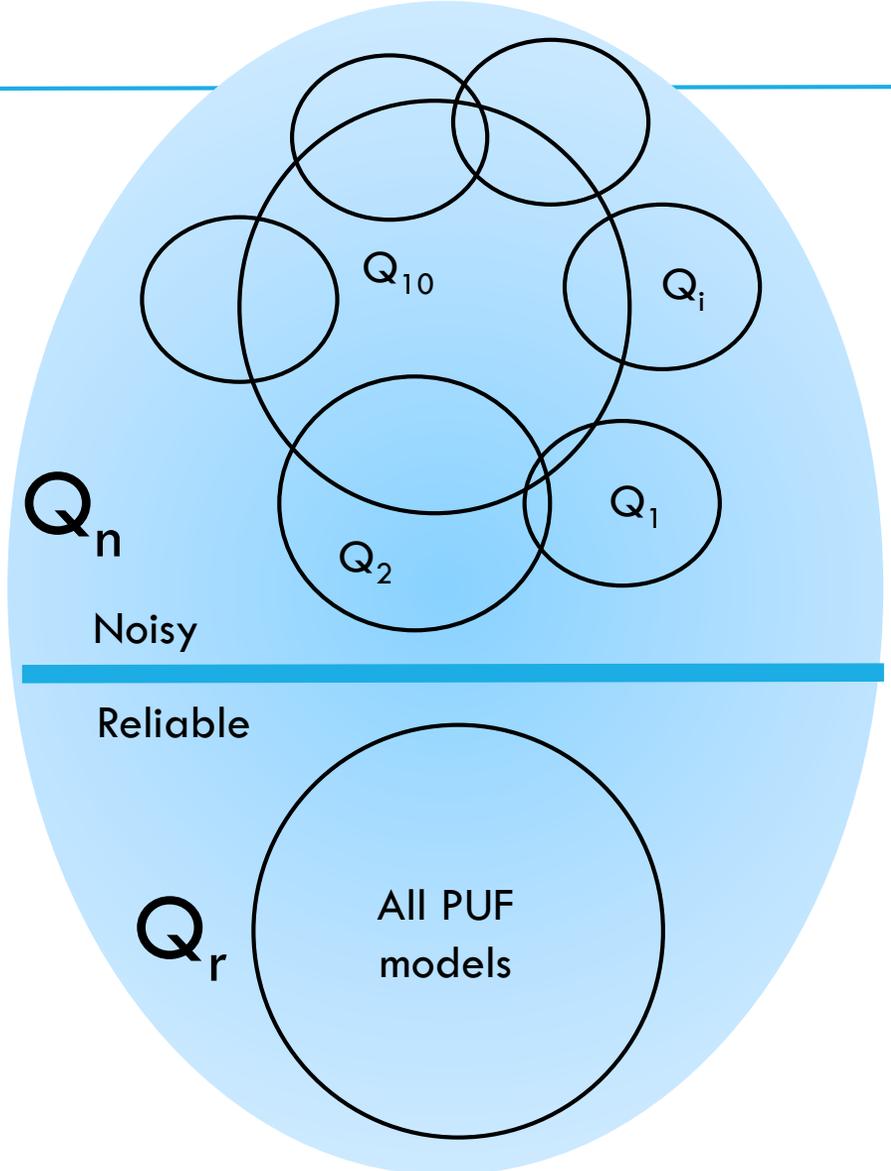
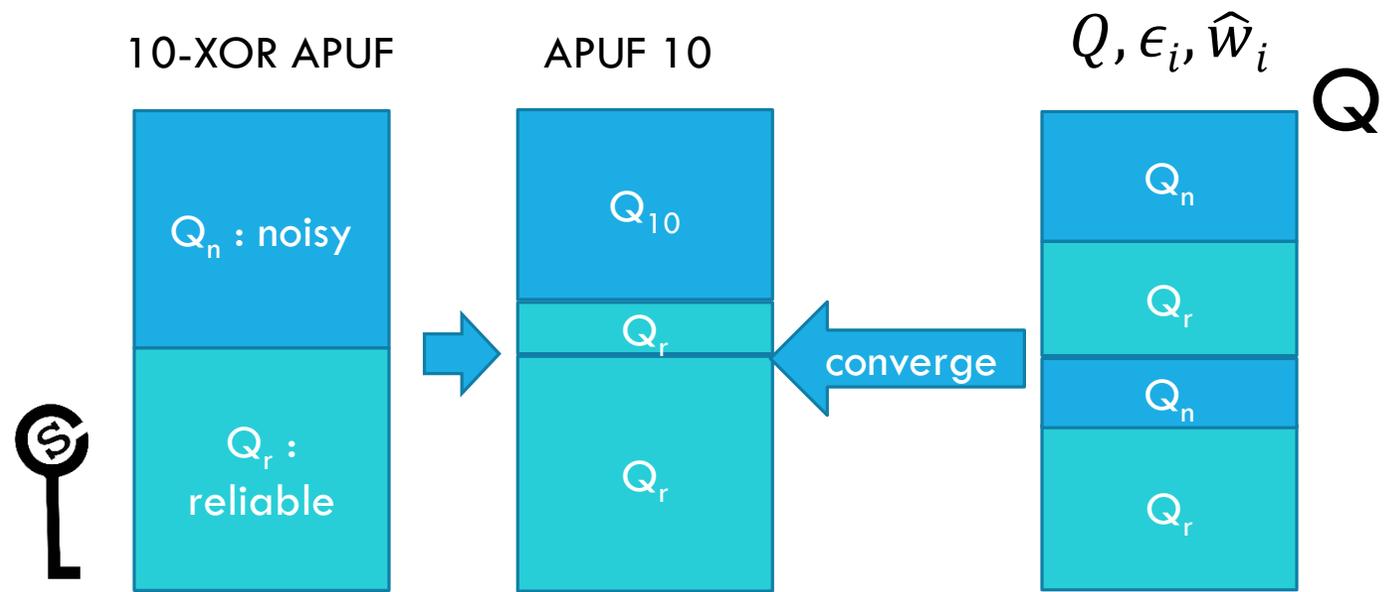
# Key idea of the attack on XOR PUF [1]

- (1) All the models  $\hat{w}_i$  in CMA-ES are **models of APUF**



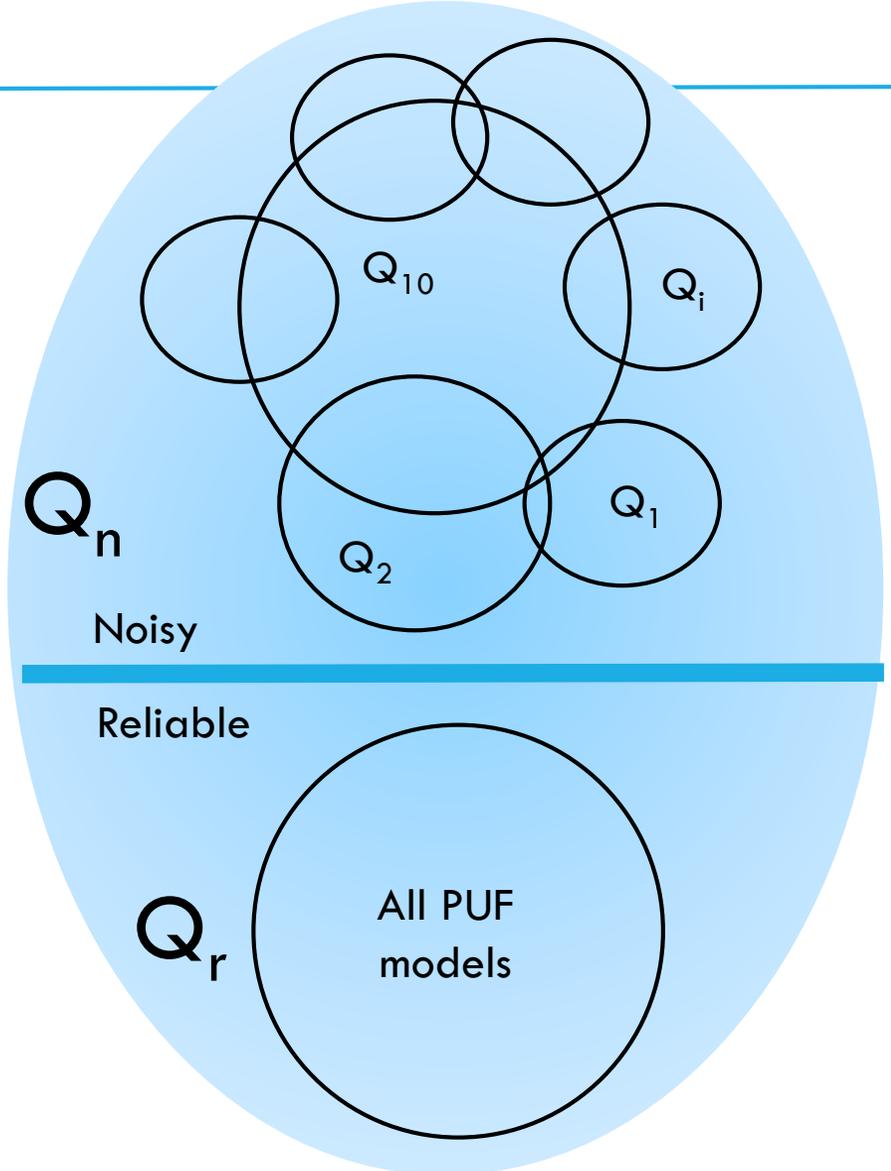
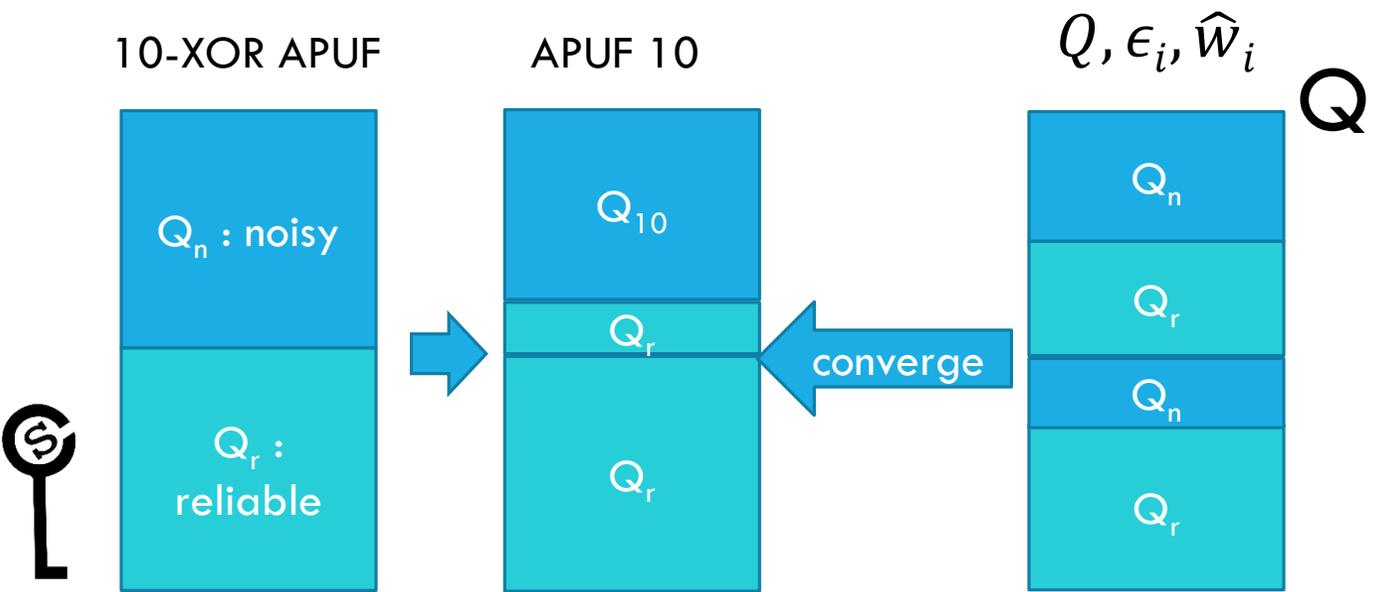
# Key idea of the attack on XOR PUF [2]

- (1) All the models  $\hat{w}_i$  in CMA-ES are **models of APUF**
- (2)  $\hat{w}_i$  can **only** converge to an APUF instance



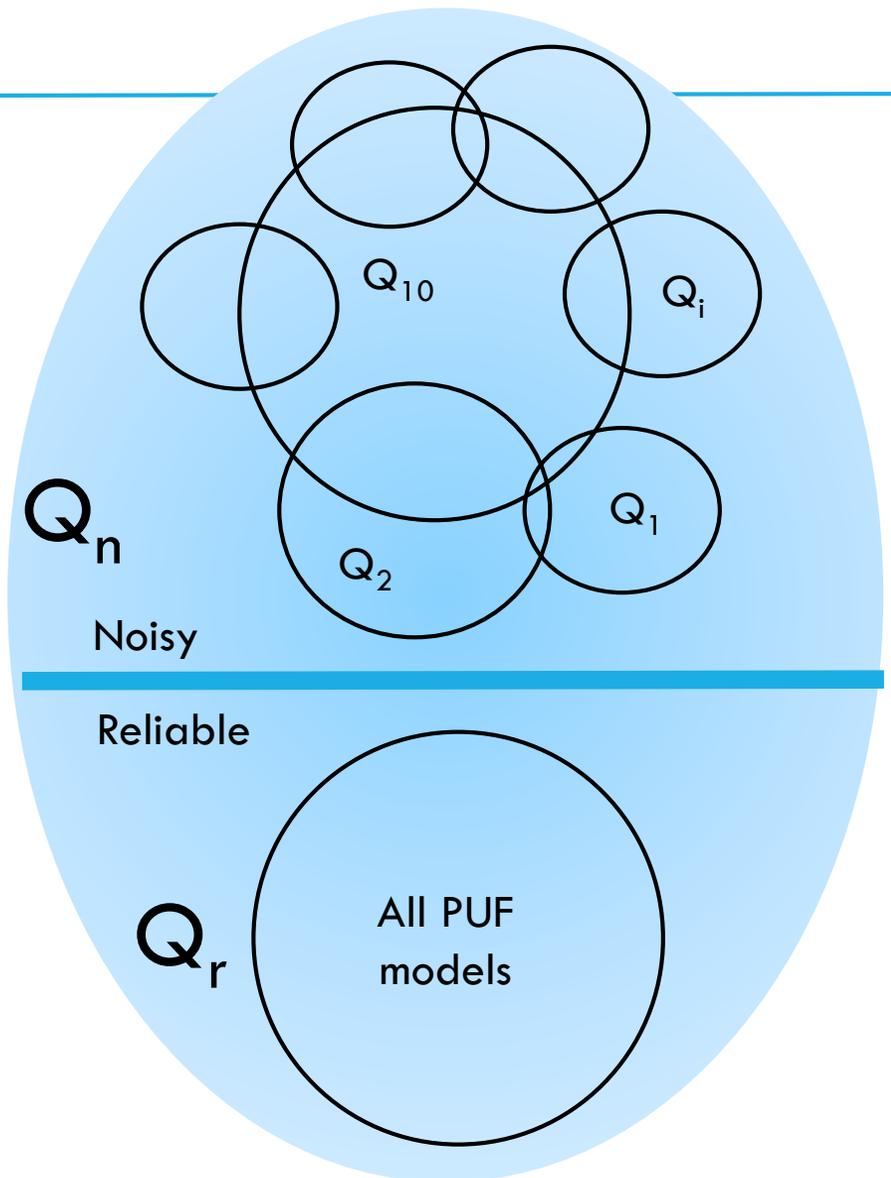
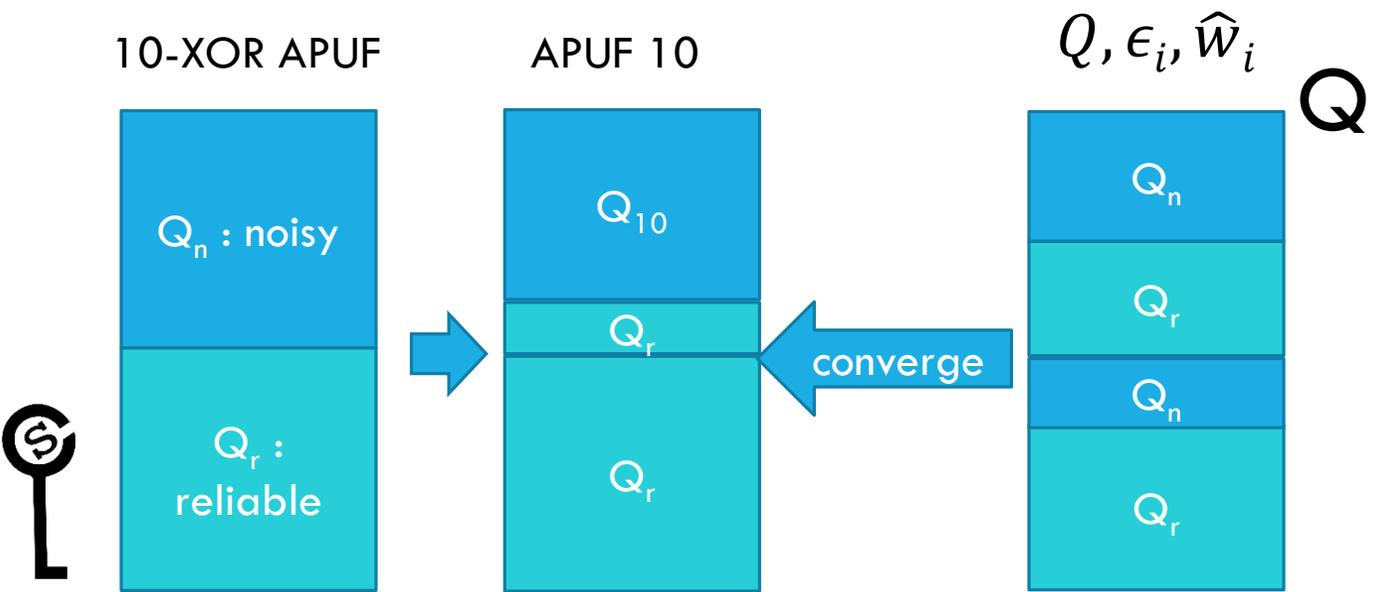
# Key idea of the attack on XOR PUF [3]

- (1) All the models  $\hat{w}_i$  in CMA-ES are **models of APUF**
- (2)  $\hat{w}_i$  can **only** converge to an APUF instance
- (3) CMA ES **maximizes the matching**  $Q$  of  $\hat{w}_i$  and  $Q$  of XOR APUF



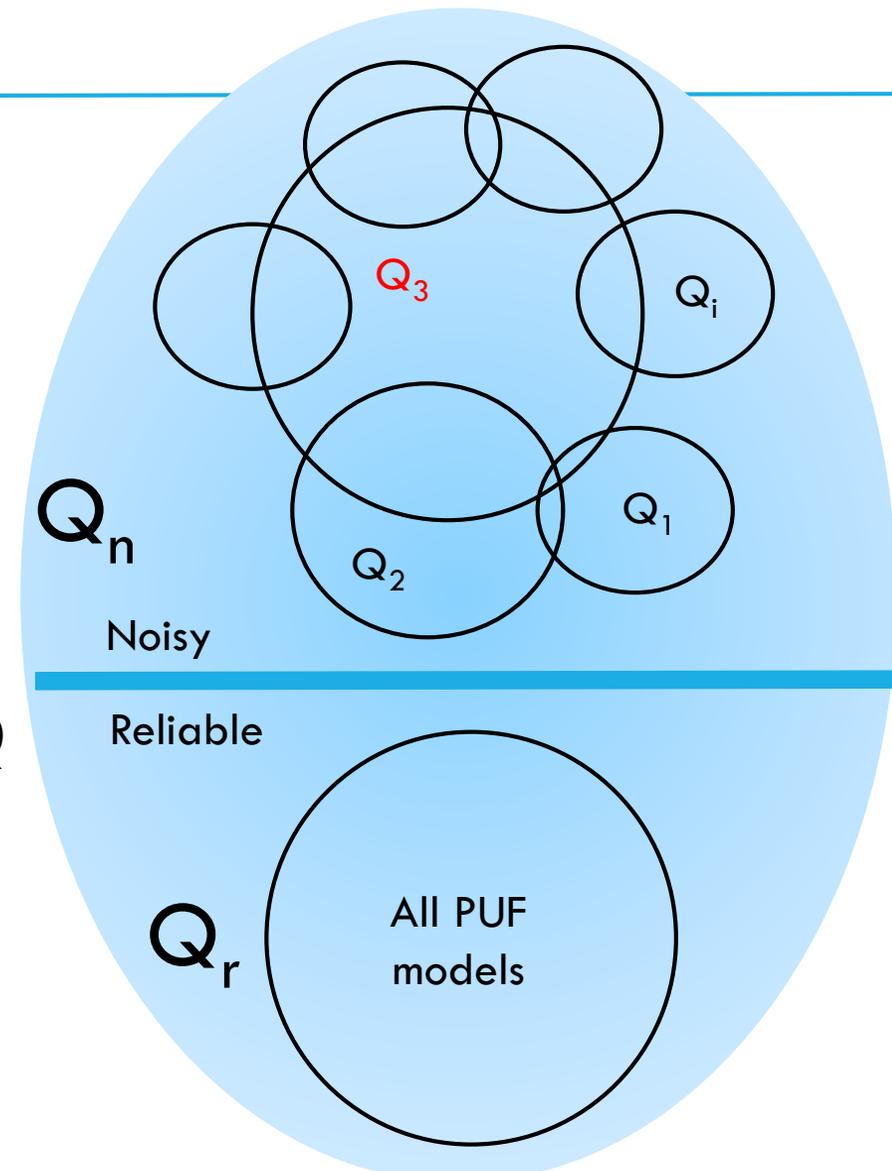
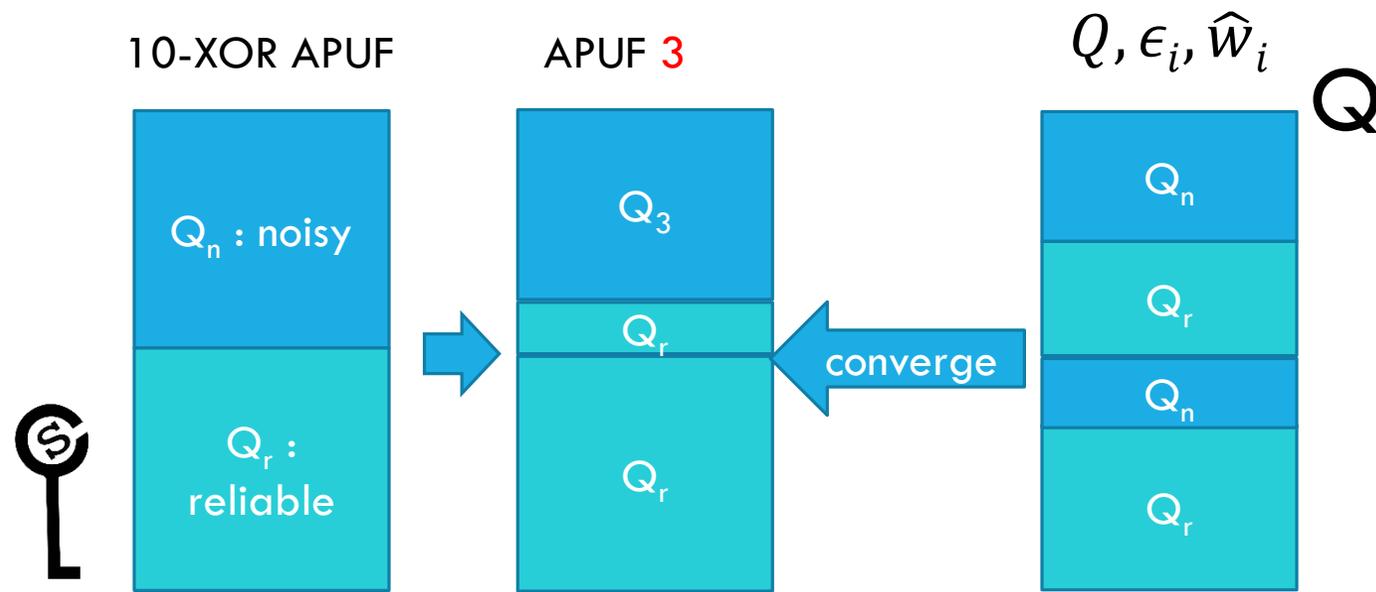
# Key idea of the attack on XOR PUF [4]

- (1) All the models  $\hat{w}_i$  in CMA-ES are **models of APUF**
- (2)  $\hat{w}_i$  can **only** converge to an APUF instance
- (3) CMA ES **maximizes the matching**  $Q$  of  $\hat{w}_i$  and  $Q$  of XOR APUF
- (1)+(2)+(3) CMA ES forces  $\hat{w}_i$  converges to APUF 10 because  **$Q$  of APUF 10 is the representative of  $Q$  of XOR APUF.**



# Key idea of the attack on XOR PUF [5]

- Changing  $Q$  makes  $Q_3$  largest



# Key idea of the attack on XOR PUF [6]

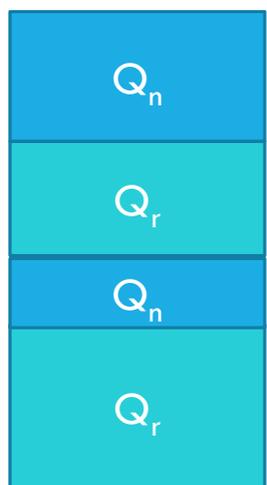
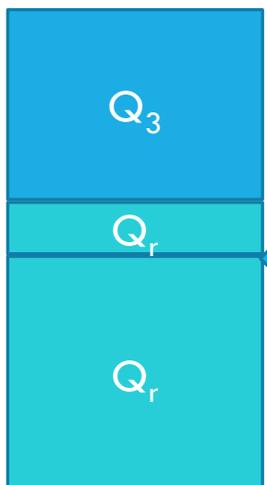
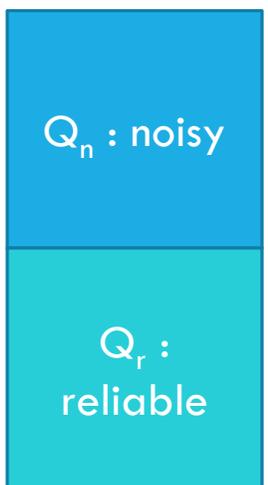
- Changing  $Q$  makes  $Q_3$  largest

Keep changing  $Q$  and applying CMA-ES attack on  $Q$  to get models of all APUF instances

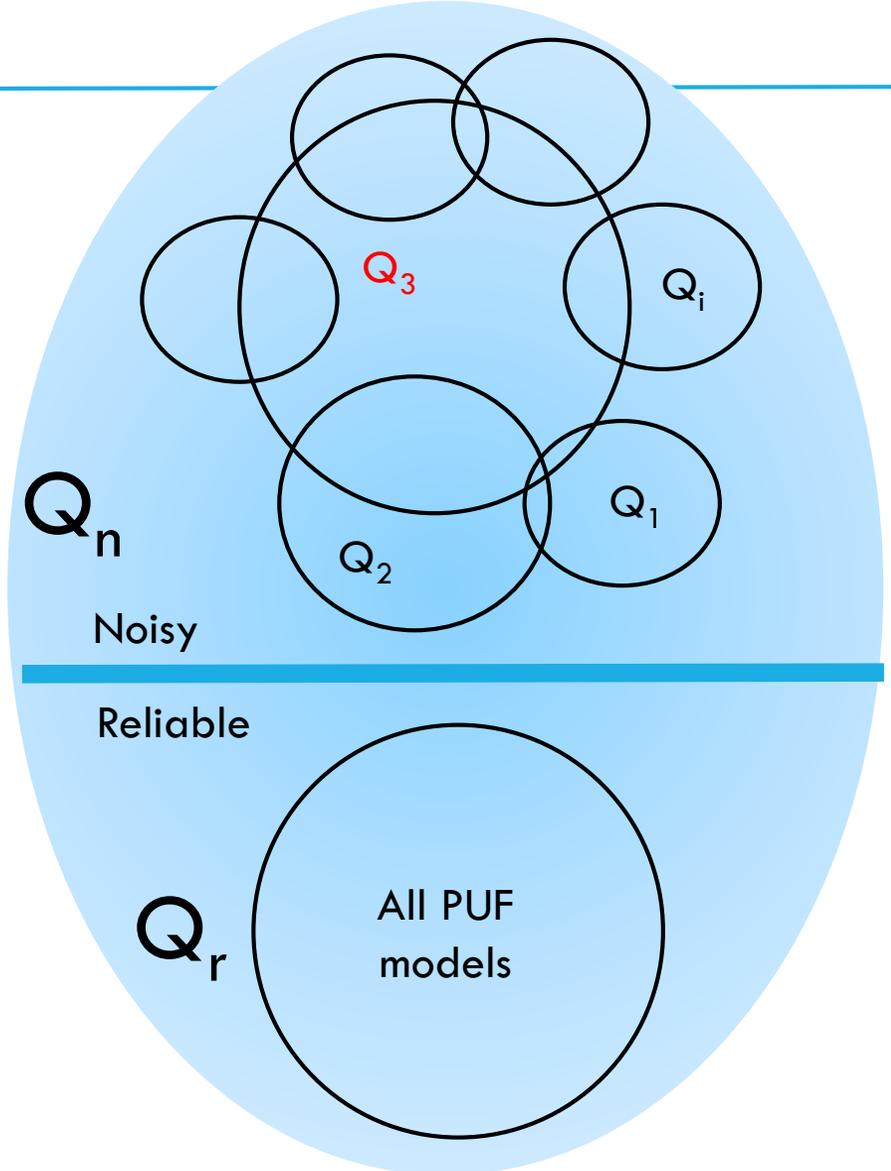
10-XOR APUF

APUF 3

$Q, \epsilon_i, \hat{W}_i$



$Q$

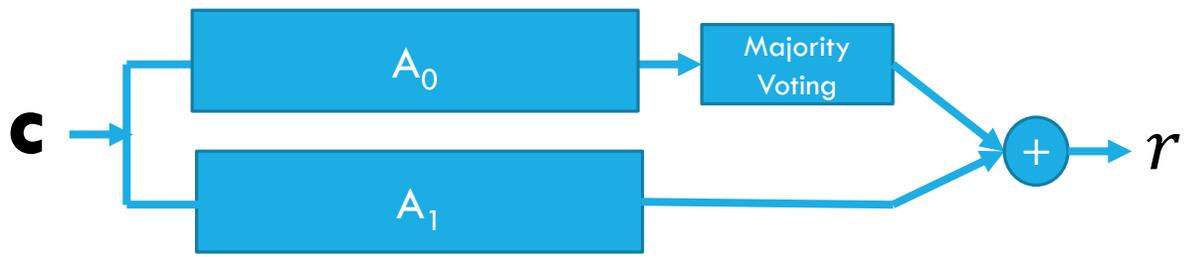


## Question 2: How to make the attack on XOR PUF fail?

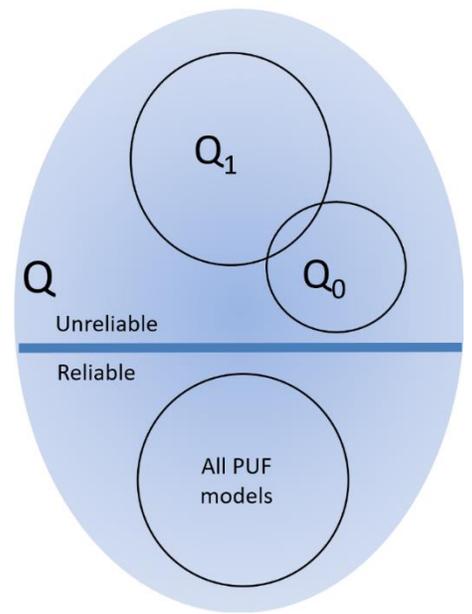
---



# Attack fails



2-XOR APUF with majority voting circuit at  $A_0$



**CMA ES never converges to APUF  $A_0$  and always converges to APUF  $A_1$  when majority voting mechanism in use.**

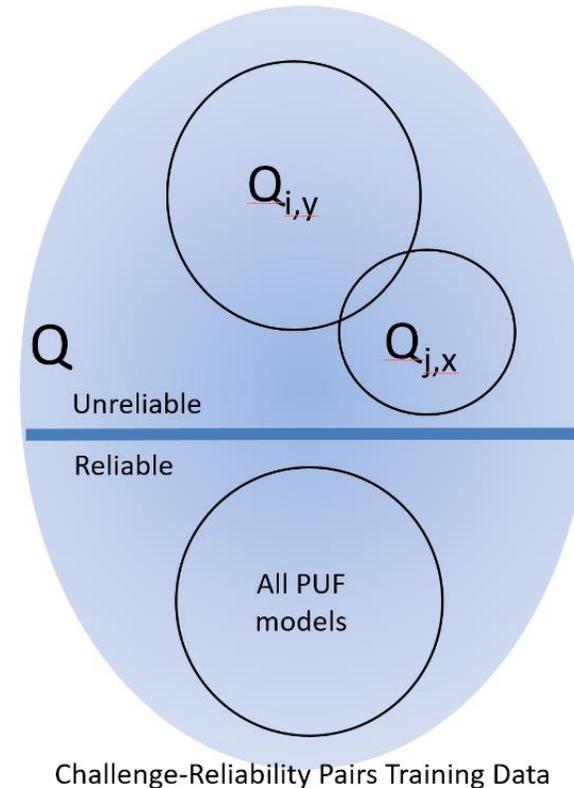
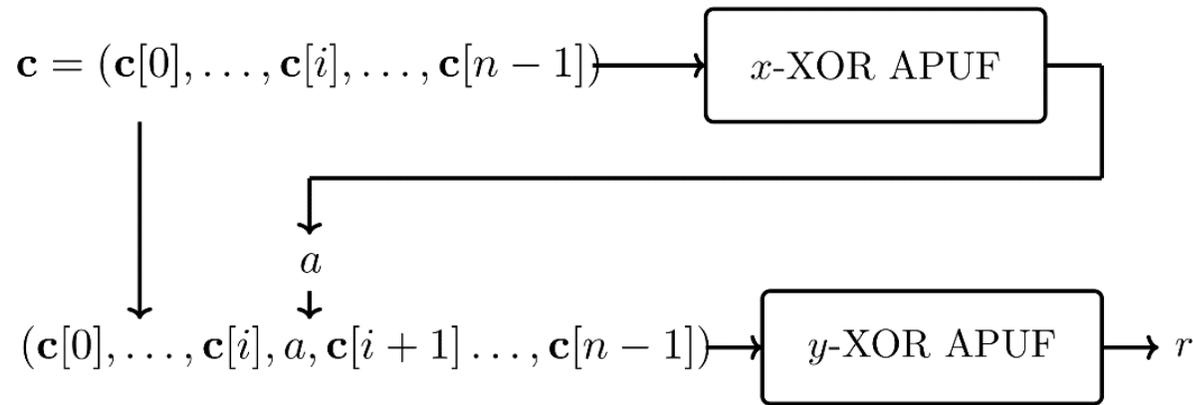


## 5. Interpose PUF (iPUF) – Reliability based modeling attack resistance



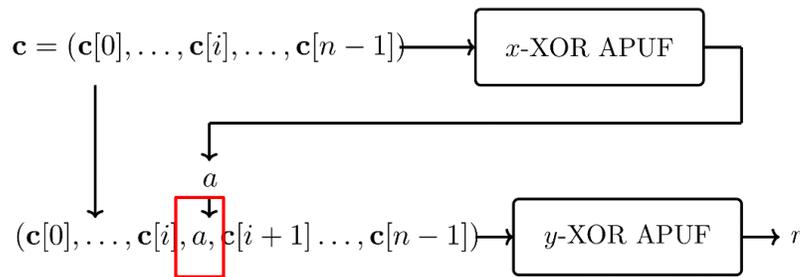
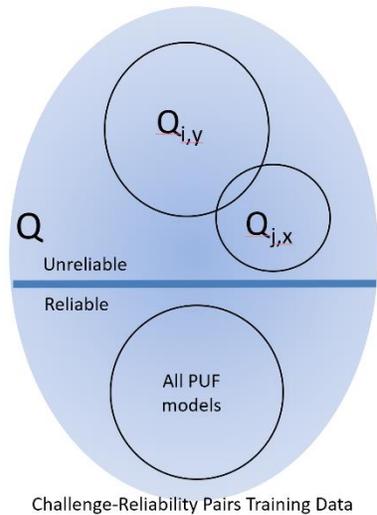
# Security of iPUF wrt Reliability-based modeling attack [1]

- Reason 1: the information of APUF instances in x-XOR PUF presented at the iPUF output is less compared to APUF instances in y-XOR PUF. Thus, the reliability based modeling attack never converges to any APUF instance in x-XOR PUF



# Security of iPUF wrt Reliability-based modeling attack [2]

- **Reason 2:** to attack APUFs at  $y$ -XOR APUF, the adversary needs to compute  $\Delta$ . But compute  $\Delta$  is infeasible because the output of  $x$ -XOR PUF ( $a$ ) is not known.



$Q, \epsilon_1, \hat{w}_1$  Cannot compute  $\Phi(c)$  or  $\Delta$

$Q \rightarrow \text{challenge } c \rightarrow \Phi(c) \rightarrow \Delta = \hat{w}_1 \cdot \Phi(c)$

$|\Delta| \leq \epsilon_1 \rightarrow \text{challenge } c \text{ is noisy} \rightarrow Q_n : \text{noisy}$

$|\Delta| > \epsilon_1 \rightarrow \text{challenge } c \text{ is reliable} \rightarrow Q_r : \text{reliable}$

$$\Delta = \Delta(n - 1) = \mathbf{w} \cdot \Phi^T$$

$$\Phi[i] = \prod_{j=i, \dots, n-1} (1 - c[j]), i = 0, \dots, n - 1$$

$$\Phi[n] = 1$$

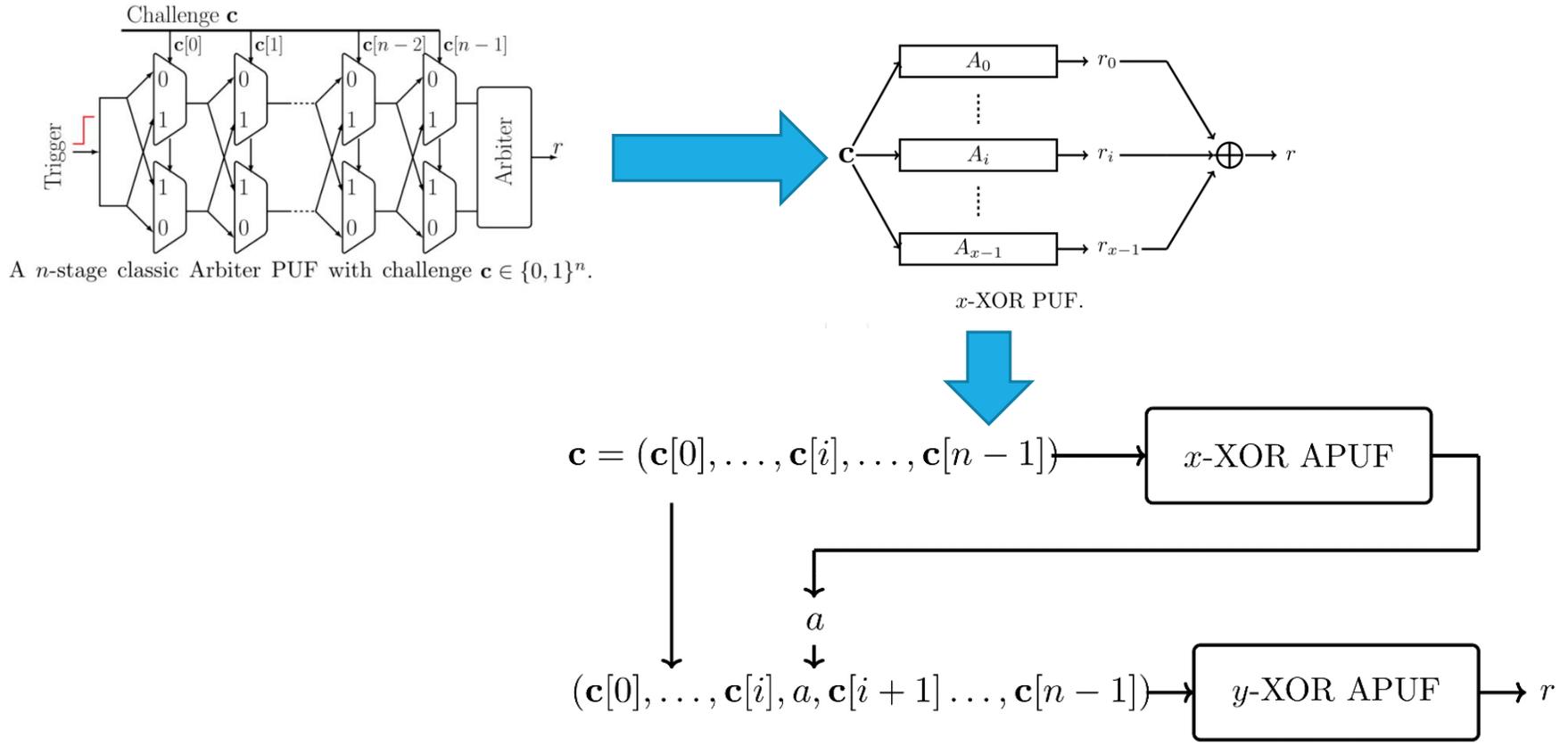


- Theoretical
  - Enhanced Reliability based Modeling Attacks on APUF and XOR APUFs
  - Proved Logistic Regression on XOR APUF is the best attack
  - Proved Logistic Regression on iPUF is not applicable
- Engineering
  - Implemented APUF, XOR, and iPUF on FPGA
  - Studied good and bad FPGA-implemented APUF based PUF
  - All source codes available online: [https://github.com/scluconn/DA\\_PUF\\_Library/](https://github.com/scluconn/DA_PUF_Library/)
- Detailed tutorial online:  
<https://www.youtube.com/playlist?list=PLK5NNs4GceLQw7bOEHSdZOwHImSF1zvS>  
[W](#)



# 6. Conclusion

- We explain how the reliability-based modeling attack on XOR PUF works
- We propose a new lightweight PUF design (iPUF) which is secure against the state-of-the-art of modelling attacks.



1. <https://slideplayer.com/slide/3927633/>
2. Cryptanalysis of electrical PUFs via machine learning algorithms – Master Thesis of Jan Solter
3. The Gap Between Promise and Reality: On the Insecurity of XOR Arbiter PUFs CHES, September 16 th , 2015, Georg T. Becker
4. <https://en.wikipedia.org/wiki/CMA-ES>

Thank you for your attention!  
and any questions?

